

The Fundamental Properties of Natural Numbers

Grzegorz Bancerek¹
Warsaw University
Białystok

Summary. Some fundamental properties of addition, multiplication, order relations, exact division, the remainder, divisibility, the least common multiple, the greatest common divisor are presented. A proof of Euclid algorithm is also given.

The article [1] provides the terminology and notation for this paper. For simplicity we adopt the following convention: x will denote an object of the type `Real`; k, l, m, n will denote objects of the type `Nat`; X will denote an object of the type `set of Real`. One can prove the following propositions:

- (1) x is Nat **implies** $x + 1$ is Nat,
- (2) **for** X **st** $0 \in X$ **&** **for** x **st** $x \in X$ **holds** $x + 1 \in X$ **for** k **holds** $k \in X$,
- (3) $k + n = n + k$,
- (4) $k + m + n = k + (m + n)$,
- (5) $k + 0 = k$ **&** $0 + k = k$,
- (6) $k \cdot n = n \cdot k$,
- (7) $k \cdot (m \cdot n) = (k \cdot m) \cdot n$,
- (8) $k \cdot 1 = k$ **&** $1 \cdot k = k$,
- (9) $k \cdot (n + m) = k \cdot n + k \cdot m$ **&** $(n + m) \cdot k = n \cdot k + m \cdot k$,
- (10) $k + m = n + m$ **or** $k + m = m + n$ **or** $m + k = m + n$ **implies** $k = n$,

¹Supported by RPBP III.24 C1

$$(11) \quad k \cdot 0 = 0 \ \& \ 0 \cdot k = 0.$$

Let us consider n, k . Let us note that it makes sense to consider the following functor on a restricted area. Then

$$n + k \quad \text{is} \quad \text{Nat}.$$

The scheme *Ind* deals with a unary predicate \mathcal{P} states that the following holds

$$\mathbf{for} \ k \ \mathbf{holds} \ \mathcal{P}[k]$$

provided the parameter satisfies the following conditions:

- $\mathcal{P}[0],$
- $\mathbf{for} \ k \ \mathbf{st} \ \mathcal{P}[k] \ \mathbf{holds} \ \mathcal{P}[k + 1].$

Let us consider n, k . Let us note that it makes sense to consider the following functor on a restricted area. Then

$$n \cdot k \quad \text{is} \quad \text{Nat}.$$

One can prove the following propositions:

$$(12) \quad k \leq n \ \& \ n \leq k \ \mathbf{implies} \ k = n,$$

$$(13) \quad k \leq n \ \& \ n \leq m \ \mathbf{implies} \ k \leq m,$$

$$(14) \quad k \leq n \ \mathbf{or} \ n \leq k,$$

$$(15) \quad k \leq k,$$

$$(16) \quad k \leq n \ \mathbf{implies}$$

$$k + m \leq n + m \ \& \ k + m \leq m + n \ \& \ m + k \leq m + n \ \& \ m + k \leq n + m,$$

$$(17) \quad k + m \leq n + m \ \mathbf{or} \ k + m \leq m + n \ \mathbf{or} \ m + k \leq m + n \ \mathbf{or} \ m + k \leq n + m \\ \mathbf{implies} \ k \leq n,$$

$$(18) \quad \mathbf{for} \ k \ \mathbf{holds} \ 0 \leq k,$$

$$(19) \quad 0 \neq k \ \mathbf{implies} \ 0 < k,$$

$$(20) \quad k \leq n \ \mathbf{implies} \ k \cdot m \leq n \cdot m \ \& \ k \cdot m \leq m \cdot n \ \& \ m \cdot k \leq n \cdot m \ \& \ m \cdot k \leq m \cdot n,$$

$$(21) \quad 0 \neq k + 1,$$

$$(22) \quad k = 0 \ \mathbf{or} \ \mathbf{ex} \ n \ \mathbf{st} \ k = n + 1,$$

$$(23) \quad k + n = 0 \ \mathbf{implies} \ k = 0 \ \& \ n = 0,$$

$$(24) \quad k \neq 0 \ \& \ (n \cdot k = m \cdot k \ \mathbf{or} \ n \cdot k = k \cdot m \ \mathbf{or} \ k \cdot n = k \cdot m) \ \mathbf{implies} \ n = m,$$

$$(25) \quad k \cdot n = 0 \text{ implies } k = 0 \text{ or } n = 0.$$

The scheme *Def_by_Ind* concerns a constant \mathcal{A} that has the type Nat , a binary functor \mathcal{F} yielding values of the type Nat and a binary predicate \mathcal{P} and states that the following holds

$$(\text{for } k \text{ ex } n \text{ st } \mathcal{P}[k, n]) \ \& \ \text{for } k, n, m \text{ st } \mathcal{P}[k, n] \ \& \ \mathcal{P}[k, m] \ \text{holds } n = m$$

provided the parameters satisfy the following condition:

- $\text{for } k, n \text{ holds}$
 $\mathcal{P}[k, n] \ \text{iff } k = 0 \ \& \ n = \mathcal{A} \ \text{or } \text{ex } m, l \text{ st } k = m + 1 \ \& \ \mathcal{P}[m, l] \ \& \ n = \mathcal{F}(k, l).$

Next we state several propositions:

$$(26) \quad \text{for } k, n \text{ st } k \leq n + 1 \ \text{holds } k \leq n \ \text{or } k = n + 1,$$

$$(27) \quad \text{for } n, k \text{ st } n \leq k \ \& \ k \leq n + 1 \ \text{holds } n = k \ \text{or } k = n + 1,$$

$$(28) \quad \text{for } k, n \text{ st } k \leq n \ \text{ex } m \text{ st } n = k + m,$$

$$(29) \quad k \leq k + m,$$

$$(30) \quad k < n \ \text{iff } k \leq n \ \& \ k \neq n,$$

$$(31) \quad \text{not } k < 0.$$

Now we present three schemes. The scheme *Comp_Ind* deals with a unary predicate \mathcal{P} states that the following holds

$$\text{for } k \ \text{holds } \mathcal{P}[k]$$

provided the parameter satisfies the following condition:

- $\text{for } k \ \text{st } \text{for } n \ \text{st } n < k \ \text{holds } \mathcal{P}[n] \ \text{holds } \mathcal{P}[k].$

The scheme *Min* concerns a unary predicate \mathcal{P} states that the following holds

$$\text{ex } k \ \text{st } \mathcal{P}[k] \ \& \ \text{for } n \ \text{st } \mathcal{P}[n] \ \text{holds } k \leq n$$

provided the parameter satisfies the following condition:

- $\text{ex } k \ \text{st } \mathcal{P}[k].$

The scheme *Max* concerns a unary predicate \mathcal{P} and a constant \mathcal{A} that has the type Nat , and states that the following holds

$$\text{ex } k \ \text{st } \mathcal{P}[k] \ \& \ \text{for } n \ \text{st } \mathcal{P}[n] \ \text{holds } n \leq k$$

provided the parameters satisfy the following conditions:

- $\text{for } k \ \text{st } \mathcal{P}[k] \ \text{holds } k \leq \mathcal{A},$

- **ex k st** $\mathcal{P}[k]$.

We now state a number of propositions:

$$(32) \quad \mathbf{not} (k < n \ \& \ n < k),$$

$$(33) \quad k < n \ \& \ n < m \ \mathbf{implies} \ k < m,$$

$$(34) \quad k < n \ \mathbf{or} \ k = n \ \mathbf{or} \ n < k,$$

$$(35) \quad \mathbf{not} \ k < k,$$

$$(36) \quad k < n \ \mathbf{implies} \\ k + m < n + m \ \& \ k + m < m + n \ \& \ m + k < m + n \ \& \ m + k < n + m,$$

$$(37) \quad k \leq n \ \mathbf{implies} \ k \leq n + m,$$

$$(38) \quad k < n + 1 \ \mathbf{iff} \ k \leq n,$$

$$(39) \quad k \leq n \ \& \ n < m \ \mathbf{or} \ k < n \ \& \ n \leq m \ \mathbf{or} \ k < n \ \& \ n < m \ \mathbf{implies} \ k < m,$$

$$(40) \quad k \cdot n = 1 \ \mathbf{implies} \ k = 1 \ \& \ n = 1,$$

$$(41) \quad k + 1 \leq n \ \mathbf{iff} \ k < n.$$

The scheme *Regr* concerns a unary predicate \mathcal{P} states that the following holds

$$\mathcal{P}[0]$$

provided the parameter satisfies the following conditions:

- **ex k st** $\mathcal{P}[k]$,
- **for k st** $k \neq 0 \ \& \ \mathcal{P}[k]$ **ex n st** $n < k \ \& \ \mathcal{P}[n]$.

In the sequel $k1, t, t1$ will denote objects of the type \mathbf{Nat} . The following propositions are true:

$$(42) \quad \mathbf{for} \ m \ \mathbf{st} \ 0 < m \ \mathbf{for} \ n \ \mathbf{ex} \ k, t \ \mathbf{st} \ n = (m \cdot k) + t \ \& \ t < m,$$

$$(43) \quad \mathbf{for} \ n, m, k, k1, t, t1 \\ \mathbf{st} \ n = m \cdot k + t \ \& \ t < m \ \& \ n = m \cdot k1 + t1 \ \& \ t1 < m \ \mathbf{holds} \ k = k1 \ \& \ t = t1.$$

We now define two new functors. Let k, l have the type \mathbf{Nat} . The functor

$$k \div l,$$

yields the type \mathbf{Nat} and is defined by

$$(\mathbf{ex} \ t \ \mathbf{st} \ k = l \cdot \mathbf{it} + t \ \& \ t < l) \ \mathbf{or} \ \mathbf{it} = 0 \ \& \ l = 0.$$

The functor

$$k \bmod l,$$

yields the type Nat and is defined by

$$(\text{ext st } k = l \cdot t + \text{it} \ \& \ \text{it} < l) \text{ or } \text{it} = 0 \ \& \ l = 0.$$

Next we state four propositions:

$$(44) \quad \text{for } k, l, n \text{ being Nat} \\ \text{holds } n = k \div l \text{ iff } (\text{ext st } k = l \cdot n + t \ \& \ t < l) \text{ or } n = 0 \ \& \ l = 0,$$

$$(45) \quad \text{for } k, l, n \text{ being Nat} \\ \text{holds } n = k \bmod l \text{ iff } (\text{ext st } k = l \cdot t + n \ \& \ n < l) \text{ or } n = 0 \ \& \ l = 0,$$

$$(46) \quad \text{for } m, n \text{ st } 0 < m \text{ holds } n \bmod m < m,$$

$$(47) \quad \text{for } n, m \text{ st } 0 < m \text{ holds } n = m \cdot (n \div m) + (n \bmod m).$$

Let k, l have the type Nat . The predicate

$$k \mid l \quad \text{is defined by} \quad \text{ext st } l = k \cdot t.$$

Next we state a number of propositions:

$$(48) \quad \text{for } k, l \text{ being Nat holds } k \mid l \text{ iff ext st } l = k \cdot t,$$

$$(49) \quad \text{for } n, m \text{ holds } m \mid n \text{ iff } n = m \cdot (n \div m),$$

$$(50) \quad \text{for } n \text{ holds } n \mid n,$$

$$(51) \quad \text{for } n, m, l \text{ st } n \mid m \ \& \ m \mid l \text{ holds } n \mid l,$$

$$(52) \quad \text{for } n, m \text{ st } n \mid m \ \& \ m \mid n \text{ holds } n = m,$$

$$(53) \quad k \mid 0 \ \& \ 1 \mid k,$$

$$(54) \quad \text{for } n, m \text{ st } 0 < m \ \& \ n \mid m \text{ holds } n \leq m,$$

$$(55) \quad \text{for } n, m, l \text{ st } n \mid m \ \& \ n \mid l \text{ holds } n \mid m + l,$$

$$(56) \quad n \mid k \text{ implies } n \mid k \cdot m,$$

$$(57) \quad \text{for } n, m, l \text{ st } n \mid m \ \& \ n \mid m + l \text{ holds } n \mid l,$$

$$(58) \quad n \mid m \ \& \ n \mid k \text{ implies } n \mid m \bmod k.$$

Let us consider k, n . The functor

$$k \text{ lcm } n,$$

with values of the type Nat , is defined by

$$k \mid \mathbf{it} \ \& \ n \mid \mathbf{it} \ \& \ \mathbf{for} \ m \ \mathbf{st} \ k \mid m \ \& \ n \mid m \ \mathbf{holds} \ \mathbf{it} \mid m.$$

Next we state a proposition

$$(59) \quad \mathbf{for} \ M \ \mathbf{being} \ \text{Nat} \\ \mathbf{holds} \ M = k \ \text{lcm} \ n \ \mathbf{iff} \ k \mid M \ \& \ n \mid M \ \& \ \mathbf{for} \ m \ \mathbf{st} \ k \mid m \ \& \ n \mid m \ \mathbf{holds} \ M \mid m.$$

Let us consider k, n . The functor

$$k \ \text{gcd} \ n,$$

yields the type Nat and is defined by

$$\mathbf{it} \mid k \ \& \ \mathbf{it} \mid n \ \& \ \mathbf{for} \ m \ \mathbf{st} \ m \mid k \ \& \ m \mid n \ \mathbf{holds} \ m \mid \mathbf{it}.$$

We now state a proposition

$$(60) \quad \mathbf{for} \ M \ \mathbf{being} \ \text{Nat} \\ \mathbf{holds} \ M = k \ \text{gcd} \ n \ \mathbf{iff} \ M \mid k \ \& \ M \mid n \ \& \ \mathbf{for} \ m \ \mathbf{st} \ m \mid k \ \& \ m \mid n \ \mathbf{holds} \ m \mid M.$$

The scheme *Euklides* deals with a unary functor \mathcal{F} yielding values of the type Nat , a constant \mathcal{A} that has the type Nat and a constant \mathcal{B} that has the type Nat , and states that the following holds

$$\mathbf{ex} \ n \ \mathbf{st} \ \mathcal{F}(n) = \mathcal{A} \ \text{gcd} \ \mathcal{B} \ \& \ \mathcal{F}(n+1) = 0$$

provided the parameters satisfy the following conditions:

- $0 < \mathcal{B} \ \& \ \mathcal{B} < \mathcal{A},$
- $\mathcal{F}(0) = \mathcal{A} \ \& \ \mathcal{F}(1) = \mathcal{B},$
- $\mathbf{for} \ n \ \mathbf{holds} \ \mathcal{F}(n+2) = \mathcal{F}(n) \ \text{mod} \ \mathcal{F}(n+1).$

References

- [1] Krzysztof Hryniewiecki. Basic properties of real numbers. *Formalized Mathematics*, 1, 1990.

Received January 11, 1989
