JOURNAL OF FORMALIZED MATHEMATICS
Volume 5, Released 1993, Published 2003
Inst. of Computer Science, Univ. of Białystok

Two Programs for SCM. Part II - Programs[1]

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Summary. We prove the correctness of two short programs for the **SCM** machine: one computes Fibonacci numbers and the other computes the *fusc* function of Dijkstra [5]. The formal definitions of these functions can be found in [4]. We prove the total correctness of the programs in two ways: by conducting inductions on computations and inductions on input data. In addition we characterize the concrete complexity of the programs as defined in [3].

MML Identifier: FIB_FUSC.

WWW: http://mizar.org/JFM/Vol5/fib_fusc.html

The articles [9], [1], [7], [2], [6], [3], [8], and [4] provide the notation and terminology for this paper.

The program computing Fib is a finite sequence of elements of the instructions of **SCM** and is defined by:

(Def. 1) The program computing Fib $= \langle \mathbf{if}\ \mathbf{d}_1 > 0\ \mathbf{goto}\ \mathbf{i}_2 \rangle \frown \langle \mathbf{halt}_{\mathbf{SCM}} \rangle \frown \langle \mathbf{d}_3{:=}\mathbf{d}_0 \rangle \frown \langle \mathrm{SubFrom}(\mathbf{d}_1, \mathbf{d}_0) \rangle \frown \langle \mathbf{if}\ \mathbf{d}_1 = 0\ \mathbf{goto}\ \mathbf{i}_1 \rangle \frown \langle \mathbf{d}_4{:=}\mathbf{d}_2 \rangle \frown \langle \mathbf{d}_2{:=}\mathbf{d}_3 \rangle \frown \langle \mathrm{AddTo}(\mathbf{d}_3, \mathbf{d}_4) \rangle \frown \langle \mathrm{goto}\ (\mathbf{i}_3) \rangle$.

We now state the proposition

(1) Let N be a natural number and s be a state with instruction counter on 0, with the program computing Fib located from 0, and $\langle 1 \rangle \frown \langle N \rangle \frown \langle 0 \rangle \frown \langle 0 \rangle$ from 0. Then

(i) s is halting,

(ii) if $N = 0$, then the complexity of $s = 1$,

(iii) if $N > 0$, then the complexity of $s = 6 \cdot N - 2$, and

(iv) $(\mathrm{Result}(s))(\mathbf{d}_3) = \mathrm{Fib}(N)$.

Let i be an integer. The functor $\mathrm{Fusc}'(i)$ yields a natural number and is defined as follows:

(Def. 2) There exists a natural number n such that $i = n$ and $\mathrm{Fusc}'(i) = \mathrm{Fusc}(n)$ or i is not a natural number and $\mathrm{Fusc}'(i) = 0$.

The program computing Fusc is a finite sequence of elements of the instructions of **SCM** and is defined by:

(Def. 3) The program computing Fusc $= \langle \mathbf{if}\ \mathbf{d}_1 = 0\ \mathbf{goto}\ \mathbf{i}_8 \rangle \frown \langle \mathbf{d}_4{:=}\mathbf{d}_0 \rangle \frown \langle \mathrm{Divide}(\mathbf{d}_1, \mathbf{d}_4) \rangle \frown \langle \mathbf{if}\ \mathbf{d}_4 = 0\ \mathbf{goto}\ \mathbf{i}_6 \rangle \frown \langle \mathrm{AddTo}(\mathbf{d}_3, \mathbf{d}_2) \rangle \frown \langle \mathrm{goto}\ (\mathbf{i}_0) \rangle \frown \langle \mathrm{AddTo}(\mathbf{d}_2, \mathbf{d}_3) \rangle \frown \langle \mathrm{goto}\ (\mathbf{i}_0) \rangle \frown \langle \mathbf{halt}_{\mathbf{SCM}} \rangle$.

[1] This work was partially supported by NSERC Grant OGP9207 while the first author visited University of Alberta, May–June 1993.

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The following propositions are true:

(2) Let N be a natural number. Suppose $N > 0$. Let s be a state with instruction counter on 0, with the program computing Fusc located from 0, and $\langle 2\rangle \frown \langle N\rangle \frown \langle 1\rangle \frown \langle 0\rangle$ from 0. Then s is halting and $(\mathrm{Result}(s))(\mathbf{d}_3) = \mathrm{Fusc}(N)$ and the complexity of $s = 6 \cdot (\lfloor \log_2 N \rfloor + 1) + 1$.

(3) Let N be a natural number, k, F_1, F_2 be natural numbers, and s be a state with instruction counter on 3, with the program computing Fib located from 0, and $\langle 1\rangle \frown \langle N\rangle \frown \langle F_1\rangle \frown \langle F_2\rangle$ from 0. Suppose $N > 0$ and $F_1 = \mathrm{Fib}(k)$ and $F_2 = \mathrm{Fib}(k+1)$. Then

(i) s is halting,

(ii) the complexity of $s = 6 \cdot N - 4$, and

(iii) there exists a natural number m such that $m = (k+N) - 1$ and $(\mathrm{Result}(s))(\mathbf{d}_2) = \mathrm{Fib}(m)$ and $(\mathrm{Result}(s))(\mathbf{d}_3) = \mathrm{Fib}(m+1)$.

(5)[1] Let n be a natural number, N, A, B be natural numbers, and s be a state with instruction counter on 0, with the program computing Fusc located from 0, and $\langle 2\rangle \frown \langle n\rangle \frown \langle A\rangle \frown \langle B\rangle$ from 0. Suppose $N > 0$ and $\mathrm{Fusc}(N) = A \cdot \mathrm{Fusc}(n) + B \cdot \mathrm{Fusc}(n+1)$. Then

(i) s is halting,

(ii) $(\mathrm{Result}(s))(\mathbf{d}_3) = \mathrm{Fusc}(N)$,

(iii) if $n = 0$, then the complexity of $s = 1$, and

(iv) if $n > 0$, then the complexity of $s = 6 \cdot (\lfloor \log_2 n \rfloor + 1) + 1$.

(6) Let N be a natural number. Suppose $N > 0$. Let s be a state with instruction counter on 0, with the program computing Fusc located from 0, and $\langle 2\rangle \frown \langle N\rangle \frown \langle 1\rangle \frown \langle 0\rangle$ from 0. Then

(i) s is halting,

(ii) $(\mathrm{Result}(s))(\mathbf{d}_3) = \mathrm{Fusc}(N)$,

(iii) if $N = 0$, then the complexity of $s = 1$, and

(iv) if $N > 0$, then the complexity of $s = 6 \cdot (\lfloor \log_2 N \rfloor + 1) + 1$.

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[1] The proposition (4) has been removed.

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Received October 8, 1993

Published January 2, 2004