# Subtrees<sup>1</sup>

## Grzegorz Bancerek Institute of Mathematics Polish Academy of Sciences

**Summary.** The concepts of root tree, the set of successors of a node in decorated tree and sets of subtrees are introduced.

MML Identifier: TREES\_9.

WWW: http://mizar.org/JFM/Vol6/trees\_9.html

The articles [15], [12], [17], [16], [2], [18], [10], [11], [8], [14], [13], [4], [1], [3], [5], [6], [7], and [9] provide the notation and terminology for this paper.

1. ROOT TREE AND SUCCESSORS OF NODE IN DECORATED TREE

Let us mention that every tree which is finite is also finite-order.

Next we state three propositions:

- (1) For every decorated tree t holds  $t \upharpoonright \varepsilon_{\mathbb{N}} = t$ .
- (2) For every tree t and for all finite sequences p, q of elements of  $\mathbb{N}$  such that  $p \cap q \in t$  holds  $t \upharpoonright (p \cap q) = t \upharpoonright p \upharpoonright q$ .
- (3) Let t be a decorated tree and p, q be finite sequences of elements of  $\mathbb{N}$ . If  $p \cap q \in \text{dom } t$ , then  $t \upharpoonright (p \cap q) = t \upharpoonright p \upharpoonright q$ .

Let  $I_1$  be a decorated tree. We say that  $I_1$  is root if and only if:

(Def. 1)  $dom I_1 = the elementary tree of 0.$ 

One can check that every decorated tree which is root is also finite. We now state three propositions:

- (4) For every decorated tree t holds t is root iff  $\emptyset \in \text{Leaves}(\text{dom } t)$ .
- (5) For every tree t and for every element p of t holds t 
  mid p = the elementary tree of 0 iff  $p \in \text{Leaves}(t)$ .
- (6) For every decorated tree t and for every node p of t holds  $t \mid p$  is root iff  $p \in \text{Leaves}(\text{dom } t)$ .

Let us observe that there exists a decorated tree which is root and there exists a decorated tree which is finite and non root.

Let x be a set. One can verify that the root tree of x is finite and root.

Let  $I_1$  be a tree. We say that  $I_1$  is finite-branching if and only if:

<sup>&</sup>lt;sup>1</sup>This article has been worked out during the visit of the author in Nagano in Summer 1994.

(Def. 2) For every element x of  $I_1$  holds succ x is finite.

Let us observe that every tree which is finite-order is also finite-branching.

Let us note that there exists a tree which is finite.

Let  $I_1$  be a decorated tree. We say that  $I_1$  is finite-order if and only if:

(Def. 3)  $dom I_1$  is finite-order.

We say that  $I_1$  is finite-branching if and only if:

(Def. 4) dom  $I_1$  is finite-branching.

Let us mention that every decorated tree which is finite is also finite-order and every decorated tree which is finite-order is also finite-branching.

Let us observe that there exists a decorated tree which is finite.

Let *t* be a finite-order decorated tree. Note that dom *t* is finite-order.

Let t be a finite-branching decorated tree. Observe that dom t is finite-branching.

Let t be a finite-branching tree and let p be an element of t. Note that succ p is finite.

The scheme FinOrdSet deals with a unary functor  $\mathcal F$  yielding a set and a finite set  $\mathcal A$ , and states that:

For every natural number n holds  $\mathcal{F}(n) \in \mathcal{A}$  iff  $n < \operatorname{card} \mathcal{A}$  provided the following requirements are met:

- For every set x such that  $x \in \mathcal{A}$  there exists a natural number n such that  $x = \mathcal{F}(n)$ ,
- For all natural numbers i, j such that i < j and  $\mathcal{F}(j) \in \mathcal{A}$  holds  $\mathcal{F}(i) \in \mathcal{A}$ , and
- For all natural numbers i, j such that  $\mathcal{F}(i) = \mathcal{F}(j)$  holds i = j.

Let X be a set. One can verify that there exists a finite sequence of elements of X which is one-to-one and empty.

The following proposition is true

(7) Let t be a finite-branching tree, p be an element of t, and n be a natural number. Then  $p \cap \langle n \rangle \in \operatorname{succ} p$  if and only if  $n < \operatorname{card} \operatorname{succ} p$ .

Let t be a finite-branching tree and let p be an element of t. The functor Succ p yields an one-to-one finite sequence of elements of t and is defined by:

(Def. 5) len Succ  $p = \operatorname{card}\operatorname{succ} p$  and rng Succ  $p = \operatorname{succ} p$  and for every natural number i such that  $i < \operatorname{len}\operatorname{Succ} p$  holds  $(\operatorname{Succ} p)(i+1) = p \ \widehat{\ }\langle i \rangle$ .

Let t be a finite-branching decorated tree and let p be a finite sequence. Let us assume that  $p \in \text{dom } t$ . The functor succ(t,p) yielding a finite sequence is defined by:

(Def. 6) There exists an element q of dom t such that q = p and  $\operatorname{succ}(t, p) = t \cdot \operatorname{Succ} q$ .

Next we state the proposition

(8) Let t be a finite-branching decorated tree. Then there exists a set x and there exists a decorated tree yielding finite sequence p such that t = x-tree(p).

Let t be a finite decorated tree and let p be a node of t. Note that  $t 
cents_p p$  is finite. Next we state the proposition

(10)<sup>1</sup> For every finite tree t and for every element p of t such that  $t = t \upharpoonright p$  holds  $p = \emptyset$ .

Let D be a non empty set and let S be a non empty subset of FinTrees(D). Observe that every element of S is finite.

<sup>&</sup>lt;sup>1</sup> The proposition (9) has been removed.

#### 2. SET OF SUBTREES OF DECORATED TREE

Let t be a decorated tree. The functor Subtrees(t) yields a set and is defined as follows:

(Def. 7) Subtrees $(t) = \{t \mid p : p \text{ ranges over nodes of } t\}$ .

Let t be a decorated tree. Observe that Subtrees(t) is constituted of decorated trees and non empty.

Let D be a non empty set and let t be a tree decorated with elements of D. Then Subtrees(t) is a non empty subset of Trees(D).

Let D be a non empty set and let t be a finite tree decorated with elements of D. Then Subtrees(t) is a non empty subset of FinTrees(D).

Let t be a finite decorated tree. One can check that every element of Subtrees(t) is finite.

In the sequel x is a set and t,  $t_1$ ,  $t_2$  are decorated trees.

Next we state four propositions:

- (11)  $x \in \text{Subtrees}(t)$  iff there exists a node n of t such that  $x = t \upharpoonright n$ .
- (12)  $t \in \text{Subtrees}(t)$ .
- (13) If  $t_1$  is finite and Subtrees $(t_1)$  = Subtrees $(t_2)$ , then  $t_1 = t_2$ .
- (14) For every node n of t holds Subtrees $(t \mid n) \subseteq \text{Subtrees}(t)$ .

Let t be a decorated tree. The functor FixedSubtrees(t) yields a subset of [: dom t, Subtrees(t):] and is defined as follows:

(Def. 8) FixedSubtrees(t) = { $\langle p, t | p \rangle$  : p ranges over nodes of t}.

Let t be a decorated tree. Observe that FixedSubtrees(t) is non empty. Next we state three propositions:

- (15)  $x \in \text{FixedSubtrees}(t)$  iff there exists a node n of t such that  $x = \langle n, t \upharpoonright n \rangle$ .
- (16)  $\langle \emptyset, t \rangle \in \text{FixedSubtrees}(t)$ .
- (17) If FixedSubtrees( $t_1$ ) = FixedSubtrees( $t_2$ ), then  $t_1 = t_2$ .

Let t be a decorated tree and let C be a set. The functor C-Subtrees(t) yields a subset of Subtrees(t) and is defined by:

(Def. 9) C-Subtrees $(t) = \{t \mid p; p \text{ ranges over nodes of } t: p \notin \text{Leaves}(\text{dom } t) \lor t(p) \in C\}.$ 

In the sequel *C* is a set.

One can prove the following propositions:

- (18)  $x \in C$ -Subtrees(t) iff there exists a node n of t such that  $x = t \upharpoonright n$  but  $n \notin Leaves(dom t)$  or  $t(n) \in C$ .
- (19) C-Subtrees(t) is empty iff t is root and  $t(\emptyset) \notin C$ .

Let t be a finite decorated tree and let C be a set. The functor C-ImmediateSubtrees(t) yielding a function from C-Subtrees(t) into (Subtrees(t))\* is defined by the condition (Def. 10).

(Def. 10) Let d be a decorated tree. Suppose  $d \in C$ -Subtrees(t). Let p be a finite sequence of elements of Subtrees(t). If p = (C-ImmediateSubtrees(t))(d), then  $d = d(\emptyset)$ -tree(p).

### 3. SET OF SUBTREES OF SET OF DECORATED TREE

Let X be a constituted of decorated trees non empty set. The functor Subtrees(X) yields a set and is defined as follows:

(Def. 11) Subtrees(X) = { $t \mid p : t$  ranges over elements of X, p ranges over nodes of t}.

Let X be a constituted of decorated trees non empty set. Note that Subtrees(X) is constituted of decorated trees and non empty.

Let D be a non empty set and let X be a non empty subset of Trees(D). Then Subtrees(X) is a non empty subset of Trees(D).

Let D be a non empty set and let X be a non empty subset of FinTrees(D). Then Subtrees(X) is a non empty subset of FinTrees(D).

In the sequel *X*, *Y* denote non empty constituted of decorated trees sets.

Next we state three propositions:

- (20)  $x \in \text{Subtrees}(X)$  iff there exists an element t of X and there exists a node n of t such that  $x = t \upharpoonright n$ .
- (21) If  $t \in X$ , then  $t \in \text{Subtrees}(X)$ .
- (22) If  $X \subseteq Y$ , then Subtrees $(X) \subseteq$  Subtrees(Y).

Let t be a decorated tree. Observe that  $\{t\}$  is non empty and constituted of decorated trees. We now state two propositions:

- (23) Subtrees( $\{t\}$ ) = Subtrees(t).
- (24) Subtrees(X) =  $\bigcup$ {Subtrees(t) : t ranges over elements of X}.

Let X be a constituted of decorated trees non empty set and let C be a set. The functor C-Subtrees(X) yielding a subset of Subtrees(X) is defined by:

(Def. 12) C-Subtrees $(X) = \{t \mid p; t \text{ ranges over elements of } X, p \text{ ranges over nodes of } t: p \notin \text{Leaves}(\text{dom } t) \lor t(p) \in C\}.$ 

One can prove the following four propositions:

- (25)  $x \in C$ -Subtrees(X) iff there exists an element t of X and there exists a node n of t such that  $x = t \upharpoonright n$  but  $n \notin \text{Leaves}(\text{dom } t)$  or  $t(n) \in C$ .
- (26) *C*-Subtrees(*X*) is empty iff for every element *t* of *X* holds *t* is root and  $t(\emptyset) \notin C$ .
- (27) C-Subtrees( $\{t\}$ ) = C-Subtrees(t).
- (28) C-Subtrees $(X) = \bigcup \{C$ -Subtrees $(t) : t \text{ ranges over elements of } X\}.$

Let X be a non empty constituted of decorated trees set. Let us assume that every element of X is finite. Let C be a set. The functor C-ImmediateSubtrees(X) yielding a function from C-Subtrees(X) into (Subtrees(X))\* is defined by the condition (Def. 13).

(Def. 13) Let d be a decorated tree. Suppose  $d \in C$ -Subtrees(X). Let p be a finite sequence of elements of Subtrees(X). If p = (C-ImmediateSubtrees(X))(d), then  $d = d(\emptyset)$ -tree(p).

Let t be a tree. Note that there exists an element of t which is empty. Next we state four propositions:

- (29) For every finite decorated tree t and for every element p of dom t holds len  $\operatorname{succ}(t, p) = \operatorname{len} \operatorname{Succ} p$  and dom  $\operatorname{succ}(t, p) = \operatorname{dom} \operatorname{Succ} p$ .
- (30) For every finite tree yielding finite sequence p and for every empty element n of p holds card succ n = len p.
- (31) Let t be a finite decorated tree, x be a set, and p be a decorated tree yielding finite sequence. Suppose t = x-tree(p). Let n be an empty element of dom t. Then succ(t, n) = t he roots of p.
- (32) For every finite decorated tree t and for every node p of t and for every node q of t 
  supp p holds  $\operatorname{succ}(t, p 
  supp q) = \operatorname{succ}(t 
  supp p, q)$ .

#### REFERENCES

- [1] Grzegorz Bancerek. Cardinal numbers. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Voll/card\_1.html.
- [2] Grzegorz Bancerek. The fundamental properties of natural numbers. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Vol1/nat\_1.html.
- [3] Grzegorz Bancerek. Introduction to trees. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Vol1/trees\_1. html.
- [4] Grzegorz Bancerek. Cartesian product of functions. Journal of Formalized Mathematics, 3, 1991. http://mizar.org/JFM/Vol3/funct\_6.html.
- [5] Grzegorz Bancerek. König's Lemma. Journal of Formalized Mathematics, 3, 1991. http://mizar.org/JFM/Vol3/trees\_2.html.
- [6] Grzegorz Bancerek. Sets and functions of trees and joining operations of trees. Journal of Formalized Mathematics, 4, 1992. http://mizar.org/JFM/Vol4/trees\_3.html.
- [7] Grzegorz Bancerek. Joining of decorated trees. Journal of Formalized Mathematics, 5, 1993. http://mizar.org/JFM/Vol5/trees\_4.html.
- [8] Grzegorz Bancerek and Krzysztof Hryniewiecki. Segments of natural numbers and finite sequences. *Journal of Formalized Mathematics*, 1, 1989. http://mizar.org/JFM/Vol1/finseq\_1.html.
- [9] Grzegorz Bancerek and Piotr Rudnicki. On defining functions on trees. Journal of Formalized Mathematics, 5, 1993. http://mizar.org/JFM/Vol5/dtconstr.html.
- [10] Czesław Byliński. Functions and their basic properties. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Voll/funct 1.html.
- [11] Czesław Byliński. Functions from a set to a set. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Vol1/funct\_
- [12] Czesław Byliński. Some basic properties of sets. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Vol1/zfmisc 1.html.
- [13] Czesław Byliński. Finite sequences and tuples of elements of a non-empty sets. Journal of Formalized Mathematics, 2, 1990. http://mizar.org/JFM/Vol2/finseq\_2.html.
- [14] Agata Darmochwał. Finite sets. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Voll/finset\_1.html.
- [15] Andrzej Trybulec. Tarski Grothendieck set theory. Journal of Formalized Mathematics, Axiomatics, 1989. http://mizar.org/JFM/Axiomatics/tarski.html.
- [16] Andrzej Trybulec. Subsets of real numbers. Journal of Formalized Mathematics, Addenda, 2003. http://mizar.org/JFM/Addenda/numbers.html.
- [17] Zinaida Trybulec. Properties of subsets. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Vol1/subset\_1.html.
- [18] Edmund Woronowicz. Relations and their basic properties. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Voll/relat\_1.html.

Received November 25, 1994

Published January 2, 2004