# **Miscellaneous Facts about Functions**

Grzegorz Bancerek Institute of Mathematics Polish Academy of Sciences Andrzej Trybulec Warsaw University Białystok

MML Identifier: FUNCT\_7.
WWW: http://mizar.org/JFM/Vol8/funct\_7.html

The articles [21], [10], [25], [23], [2], [17], [22], [20], [1], [16], [26], [7], [27], [6], [5], [15], [11], [14], [9], [19], [8], [12], [13], [3], [24], [18], and [4] provide the notation and terminology for this paper.

## 1. PRELIMINARIES

- We follow the rules: *a*, *x*, *A*, *B* denote sets and *m*, *n* denote natural numbers. The following propositions are true:
  - (1) For every function *f* and for every set *X* such that  $\operatorname{rng} f \subseteq X$  holds  $\operatorname{id}_X \cdot f = f$ .
  - (2) Let X be a set, Y be a non empty set, and f be a function from X into Y. Suppose f is one-to-one. Let B be a subset of X and C be a subset of Y. If  $C \subseteq f^{\circ}B$ , then  $f^{-1}(C) \subseteq B$ .
  - (3) Let X, Y be non empty sets and f be a function from X into Y. Suppose f is one-to-one. Let x be an element of X and A be a subset of X. If  $f(x) \in f^{\circ}A$ , then  $x \in A$ .
  - (4) Let *X*, *Y* be non empty sets and *f* be a function from *X* into *Y*. Suppose *f* is one-to-one. Let *x* be an element of *X*, *A* be a subset of *X*, and *B* be a subset of *Y*. If  $f(x) \in f^{\circ}A \setminus B$ , then  $x \in A \setminus f^{-1}(B)$ .
  - (5) Let X, Y be non empty sets and f be a function from X into Y. Suppose f is one-to-one. Let y be an element of Y, A be a subset of X, and B be a subset of Y. If  $y \in f^{\circ}A \setminus B$ , then  $f^{-1}(y) \in A \setminus f^{-1}(B)$ .
  - (6) For every function f and for every set a such that  $a \in \text{dom } f \text{ holds } f \upharpoonright \{a\} = a \mapsto f(a)$ .

Let *x*, *y* be sets. Observe that  $x \mapsto y$  is non empty. Let *x*, *y*, *a*, *b* be sets. Observe that  $[x \mapsto a, y \mapsto b]$  is non empty. One can prove the following propositions:

- (7) For every set *I* and for every many sorted set *M* indexed by *I* and for every set *i* such that  $i \in I$  holds  $i \mapsto M(i) = M \upharpoonright \{i\}$ .
- (8) Let *I*, *J* be sets, *M* be a many sorted set indexed by [:I, J:], and *i*, *j* be sets. If  $i \in I$  and  $j \in J$ , then  $[\langle i, j \rangle \mapsto M(i, j)] = M \upharpoonright [:\{i\}, \{j\}:]$ .

- (10)<sup>1</sup> For all functions f, g, h such that  $\operatorname{rng} g \subseteq \operatorname{dom} f$  and  $\operatorname{rng} h \subseteq \operatorname{dom} f$  holds  $f \cdot (g + \cdot h) = f \cdot g + \cdot f \cdot h$ .
- (11) For all functions f, g, h holds  $(g+\cdot h) \cdot f = g \cdot f + \cdot h \cdot f$ .
- (12) For all functions f, g, h such that rng f misses dom g holds  $(h+\cdot g) \cdot f = h \cdot f$ .
- (13) For all sets A, B and for every set y such that A meets  $rng(id_B + \cdot (A \mapsto y))$  holds  $y \in A$ .
- (14) For all sets x, y and for every set A such that  $x \neq y$  holds  $x \notin \operatorname{rng}(\operatorname{id}_A + (x \mapsto y))$ .
- (15) For every set X and for every set a and for every function f such that dom  $f = X \cup \{a\}$  holds  $f = f | X + (a \mapsto f(a))$ .
- (16) For every function f and for all sets X, y, z holds  $f + (X \mapsto y) + (X \mapsto z) = f + (X \mapsto z)$ .
- (17) If 0 < m and  $m \le n$ , then  $\mathbb{Z}_m \subseteq \mathbb{Z}_n$ .
- (18)  $\mathbb{Z} \neq \mathbb{Z}^*$ .
- $(19) \quad \emptyset^* = \{\emptyset\}.$
- (20)  $\langle x \rangle \in A^*$  iff  $x \in A$ .
- (21)  $A \subseteq B$  iff  $A^* \subseteq B^*$ .
- (22) For every subset A of  $\mathbb{N}$  such that for all n, m such that  $n \in A$  and m < n holds  $m \in A$  holds A is a cardinal number.
- (23) Let *A* be a finite set and *X* be a non empty family of subsets of *A*. Then there exists an element *C* of *X* such that for every element *B* of *X* such that  $B \subseteq C$  holds B = C.
- (24) Let p, q be finite sequences. Suppose len p = len q + 1. Let i be a natural number. Then  $i \in \text{dom } q$  if and only if the following conditions are satisfied:
  - (i)  $i \in \operatorname{dom} p$ , and
- (ii)  $i+1 \in \operatorname{dom} p$ .

Let us note that there exists a finite sequence which is function yielding, non empty, and nonempty.

Observe that  $\emptyset$  is function yielding. Let f be a function. One can verify that  $\langle f \rangle$  is function yielding. Let g be a function. One can verify that  $\langle f, g \rangle$  is function yielding. Let h be a function. Note that  $\langle f, g, h \rangle$  is function yielding.

Let *n* be a natural number and let *f* be a function. Note that  $n \mapsto f$  is function yielding.

Let p be a finite sequence and let q be a non empty finite sequence. Observe that  $p \cap q$  is non empty and  $q \cap p$  is non empty.

Let *p*, *q* be function yielding finite sequences. Observe that  $p \cap q$  is function yielding. One can prove the following proposition

(25) Let p, q be finite sequences. Suppose  $p \cap q$  is function yielding. Then p is function yielding and q is function yielding.

<sup>&</sup>lt;sup>1</sup> The proposition (9) has been removed.

## 2. Some useful schemes

In this article we present several logical schemes. The scheme Kappa2D deals with non empty sets  $\mathcal{A}, \mathcal{B}, \mathcal{C}$  and a binary functor  $\mathcal{F}$  yielding a set, and states that:

There exists a function f from  $[:\mathcal{A}, \mathcal{B}:]$  into  $\mathcal{C}$  such that for every element x of  $\mathcal{A}$  and

for every element *y* of  $\mathcal{B}$  holds  $f(\langle x, y \rangle) = \mathcal{F}(x, y)$ 

provided the following requirement is met:

• For every element *x* of  $\mathcal{A}$  and for every element *y* of  $\mathcal{B}$  holds  $\mathcal{F}(x, y) \in \mathcal{C}$ .

The scheme *FinMono* deals with a set  $\mathcal{A}$ , a non empty set  $\mathcal{B}$ , and two unary functors  $\mathcal{F}$  and  $\mathcal{G}$  yielding sets, and states that:

 $\{\mathcal{F}(d); d \text{ ranges over elements of } \mathcal{B}: \mathcal{G}(d) \in \mathcal{A}\}$  is finite provided the parameters have the following properties:

•  $\mathcal{A}$  is finite, and

• For all elements  $d_1$ ,  $d_2$  of  $\mathcal{B}$  such that  $\mathcal{G}(d_1) = \mathcal{G}(d_2)$  holds  $d_1 = d_2$ .

The scheme *CardMono* deals with a set  $\mathcal{A}$ , a non empty set  $\mathcal{B}$ , and a unary functor  $\mathcal{F}$  yielding a set, and states that:

 $\mathcal{A} \approx \{d; d \text{ ranges over elements of } \mathcal{B} : \mathcal{F}(d) \in \mathcal{A}\}$ 

provided the following conditions are met:

• For every set x such that  $x \in \mathcal{A}$  there exists an element d of  $\mathcal{B}$  such that  $x = \mathcal{F}(d)$ , and

• For all elements  $d_1$ ,  $d_2$  of  $\mathcal{B}$  such that  $\mathcal{F}(d_1) = \mathcal{F}(d_2)$  holds  $d_1 = d_2$ .

The scheme *CardMono'* deals with a set  $\mathcal{A}$ , a non empty set  $\mathcal{B}$ , and a unary functor  $\mathcal{F}$  yielding a set, and states that:

 $\mathcal{A} \approx \{\mathcal{F}(d); d \text{ ranges over elements of } \mathcal{B}: d \in \mathcal{A}\}$ 

- provided the following requirements are met:
  - $\mathcal{A} \subseteq \mathcal{B}$ , and
  - For all elements  $d_1$ ,  $d_2$  of  $\mathcal{B}$  such that  $\mathcal{F}(d_1) = \mathcal{F}(d_2)$  holds  $d_1 = d_2$ .

The scheme  $\mathit{FuncSeqInd}$  concerns a unary predicate  $\mathcal{P},$  and states that:

For every function yielding finite sequence p holds  $\mathcal{P}[p]$ 

provided the parameters meet the following requirements:

- $\mathcal{P}[\emptyset]$ , and
- For every function yielding finite sequence p such that P[p] and for every function f holds P[p ^ ⟨f⟩].

#### 3. SOME AUXILIARY CONCEPTS

Let *x*, *y* be sets. Let us assume that  $x \in y$ . The functor  $x (\in y)$  yielding an element of *y* is defined as follows:

(Def. 1)  $x (\in y) = x$ .

One can prove the following proposition

(26) If  $x \in A \cap B$ , then  $x \in A = x \in B$ .

Let f, g be functions and let A be a set. We say that f and g are equal outside A if and only if:

(Def. 2)  $f \upharpoonright (\operatorname{dom} f \setminus A) = g \upharpoonright (\operatorname{dom} g \setminus A).$ 

We now state several propositions:

- (27) For every function f and for every set A holds f and f are equal outside A.
- (28) Let f, g be functions and A be a set. Suppose f and g are equal outside A. Then g and f are equal outside A.
- (29) Let f, g, h be functions and A be a set. Suppose f and g are equal outside A and g and h are equal outside A. Then f and h are equal outside A.

- (30) For all functions f, g and for every set A such that f and g are equal outside A holds dom  $f \setminus A = \text{dom } g \setminus A$ .
- (31) For all functions f, g and for every set A such that dom  $g \subseteq A$  holds f and f+g are equal outside A.
- Let f be a function and let i, x be sets. The functor f + (i,x) yielding a function is defined by:

(Def. 3) 
$$f + (i,x) = \begin{cases} f + (i \mapsto x), \text{ if } i \in \text{dom } f, \\ f, \text{ otherwise.} \end{cases}$$

We now state several propositions:

- (32) For every function f and for all sets d, i holds dom(f + (i, d)) = dom f.
- (33) For every function f and for all sets d, i such that  $i \in \text{dom } f$  holds (f + (i,d))(i) = d.
- (34) For every function f and for all sets d, i, j such that  $i \neq j$  holds (f + (i,d))(j) = f(j).
- (35) For every function f and for all sets d, e, i, j such that  $i \neq j$  holds f + (i,d) + (j,e) = f + (j,e) + (i,d).
- (36) For every function f and for all sets d, e, i holds f + (i,d) + (i,e) = f + (i,e).
- (37) For every function *f* and for every set *i* holds f + (i, f(i)) = f.

Let f be a finite sequence, let i be a natural number, and let x be a set. Note that f + (i,x) is finite sequence-like.

Let *D* be a set, let *f* be a finite sequence of elements of *D*, let *i* be a natural number, and let *d* be an element of *D*. Then f + (i,d) is a finite sequence of elements of *D*.

Next we state three propositions:

- (38) Let *D* be a non empty set, *f* be a finite sequence of elements of *D*, *d* be an element of *D*, and *i* be a natural number. If  $i \in \text{dom } f$ , then  $(f + (i,d))_i = d$ .
- (39) Let *D* be a non empty set, *f* be a finite sequence of elements of *D*, *d* be an element of *D*, and *i*, *j* be natural numbers. If  $i \neq j$  and  $j \in \text{dom } f$ , then  $(f + (i,d))_i = f_i$ .
- (40) Let D be a non empty set, f be a finite sequence of elements of D, d, e be elements of D, and i be a natural number. Then  $f + (i, f_i) = f$ .
  - 4. ON THE COMPOSITION OF A FINITE SEQUENCE OF FUNCTIONS

Let *X* be a set and let *p* be a function yielding finite sequence. The functor  $compose_X p$  yields a function and is defined by the condition (Def. 4).

- (Def. 4) There exists a many sorted function f indexed by  $\mathbb{N}$  such that
  - (i) compose<sub>*X*</sub>  $p = f(\operatorname{len} p)$ ,
  - (ii)  $f(0) = id_X$ , and
  - (iii) for every natural number *i* such that  $i + 1 \in \text{dom } p$  and for all functions *g*, *h* such that g = f(i) and h = p(i+1) holds  $f(i+1) = h \cdot g$ .

Let *p* be a function yielding finite sequence and let *x* be a set. The functor apply(p,x) yields a finite sequence and is defined by the conditions (Def. 5).

(Def. 5)(i) len apply $(p,x) = \operatorname{len} p + 1$ ,

- (ii) (apply(p,x))(1) = x, and
- (iii) for every natural number *i* and for every function *f* such that  $i \in \text{dom } p$  and f = p(i) holds (apply(p,x))(i+1) = f((apply(p,x))(i)).

We adopt the following convention: X, Y, x denote sets, p, q denote function yielding finite sequences, and f, g, h denote functions.

Next we state a number of propositions:

- (41)  $\operatorname{compose}_X \emptyset = \operatorname{id}_X.$
- (42) apply( $\emptyset$ , x) =  $\langle x \rangle$ .
- (43)  $\operatorname{compose}_X(p \cap \langle f \rangle) = f \cdot \operatorname{compose}_X p.$
- (44) apply $(p \land \langle f \rangle, x) = (apply(p, x)) \land \langle f((apply(p, x))(len p + 1)) \rangle$ .
- (45)  $\operatorname{compose}_{X}(\langle f \rangle \cap p) = \operatorname{compose}_{f^{\circ}X} p \cdot (f \upharpoonright X).$
- (46) apply $(\langle f \rangle \cap p, x) = \langle x \rangle \cap apply(p, f(x)).$
- (47) compose<sub>*X*</sub>  $\langle f \rangle = f \cdot id_X$ .
- (48) If dom  $f \subseteq X$ , then compose<sub>*X*</sub> $\langle f \rangle = f$ .
- (49) apply $(\langle f \rangle, x) = \langle x, f(x) \rangle$ .
- (50) If rng compose<sub>*X*</sub>  $p \subseteq Y$ , then compose<sub>*X*</sub>  $(p \cap q) = \text{compose}_Y q \cdot \text{compose}_X p$ .
- (51)  $(\operatorname{apply}(p \cap q, x))(\operatorname{len}(p \cap q) + 1) = (\operatorname{apply}(q, (\operatorname{apply}(p, x))(\operatorname{len} p + 1)))(\operatorname{len} q + 1).$
- (52)  $\operatorname{apply}(p \cap q, x) = (\operatorname{apply}(p, x))^{\circ} \operatorname{apply}(q, (\operatorname{apply}(p, x))(\operatorname{len} p + 1)).$
- (53)  $\operatorname{compose}_X \langle f, g \rangle = g \cdot f \cdot \operatorname{id}_X.$
- (54) If dom  $f \subseteq X$  or dom $(g \cdot f) \subseteq X$ , then compose<sub>*X*</sub> $\langle f, g \rangle = g \cdot f$ .
- (55) apply( $\langle f, g \rangle, x$ ) =  $\langle x, f(x), g(f(x)) \rangle$ .
- (56)  $\operatorname{compose}_X \langle f, g, h \rangle = h \cdot g \cdot f \cdot \operatorname{id}_X.$
- (57) If dom  $f \subseteq X$  or dom $(g \cdot f) \subseteq X$  or dom $(h \cdot g \cdot f) \subseteq X$ , then compose<sub>X</sub> $\langle f, g, h \rangle = h \cdot g \cdot f$ .
- (58) apply( $\langle f, g, h \rangle, x$ ) =  $\langle x \rangle \cap \langle f(x), g(f(x)), h(g(f(x))) \rangle$ .
  - Let *F* be a finite sequence. The functor firstdom(F) is defined by:
- (Def. 6)(i) firstdom(F) is empty if F is empty,
  - (ii) firstdom(F) =  $\pi_1(F(1))$ , otherwise.

The functor lastrng(F) is defined as follows:

- (Def. 7)(i) lastrng(F) is empty if F is empty,
  - (ii) lastrng(F) =  $\pi_2(F(\text{len }F))$ , otherwise.

The following three propositions are true:

- (59) firstdom( $\emptyset$ ) =  $\emptyset$  and lastrng( $\emptyset$ ) =  $\emptyset$ .
- (60) For every finite sequence p holds firstdom $(\langle f \rangle \cap p) = \text{dom } f$  and  $\text{lastrng}(p \cap \langle f \rangle) = \text{rng } f$ .
- (61) For every function yielding finite sequence p such that  $p \neq \emptyset$  holds  $\operatorname{rng\,compose}_X p \subseteq \operatorname{lastrng}(p)$ .

Let  $I_1$  be a finite sequence. We say that  $I_1$  is composable if and only if:

(Def. 8) There exists a finite sequence p such that  $\text{len } p = \text{len } I_1 + 1$  and for every natural number i such that  $i \in \text{dom } I_1$  holds  $I_1(i) \in p(i+1)^{p(i)}$ .

One can prove the following proposition

(62) For all finite sequences p, q such that  $p \cap q$  is composable holds p is composable and q is composable.

One can check that every finite sequence which is composable is also function yielding. Let us observe that every finite sequence which is empty is also composable. Let f be a function. Observe that  $\langle f \rangle$  is composable.

Let us note that there exists a finite sequence which is composable, non empty, and non-empty. A composable sequence is a composable finite sequence. Next we state several propositions:

- (63) For every composable sequence p such that  $p \neq \emptyset$  holds dom compose<sub>x</sub>  $p = \text{firstdom}(p) \cap X$ .
- (64) For every composable sequence p holds dom compose<sub>firstdom(p)</sub> p = firstdom(p).
- (65) For every composable sequence p and for every function f such that  $\operatorname{rng} f \subseteq \operatorname{firstdom}(p)$  holds  $\langle f \rangle \cap p$  is a composable sequence.
- (66) For every composable sequence p and for every function f such that  $\text{lastrng}(p) \subseteq \text{dom } f$  holds  $p \cap \langle f \rangle$  is a composable sequence.
- (67) For every composable sequence p such that  $x \in \text{firstdom}(p)$  and  $x \in X$  holds  $(\operatorname{apply}(p, x))(\operatorname{len} p + 1) = (\operatorname{compose}_X p)(x).$

Let X, Y be sets. Let us assume that if Y is empty, then X is empty. A composable sequence is said to be a composable sequence from X into Y if:

(Def. 9) firstdom(it) = X and lastrng(it)  $\subseteq Y$ .

Let Y be a non empty set, let X be a set, and let F be a composable sequence from X into Y. Then  $compose_X F$  is a function from X into Y.

Let q be a non-empty non empty finite sequence. A finite sequence is called a composable sequence along q if:

(Def. 10) len it +1 = len q and for every natural number i such that  $i \in \text{dom it holds it}(i) \in q(i+1)^{q(i)}$ .

Let q be a non-empty non empty finite sequence. Note that every composable sequence along q is composable and non-empty.

One can prove the following two propositions:

- (68) Let q be a non-empty non empty finite sequence and p be a composable sequence along q. If  $p \neq \emptyset$ , then firstdom(p) = q(1) and lastrng $(p) \subseteq q(\ln q)$ .
- (69) Let q be a non-empty non empty finite sequence and p be a composable sequence along q. Then dom compose<sub>q(1)</sub> p = q(1) and  $\operatorname{rng\,compose}_{q(1)} p \subseteq q(\operatorname{len} q)$ .

Let *f* be a function and let *n* be an element of  $\mathbb{N}$ . The functor  $f^n$  yielding a function is defined by the condition (Def. 11).

(Def. 11) There exists a function p from  $\mathbb{N}$  into  $(\operatorname{dom} f \cup \operatorname{rng} f) \rightarrow (\operatorname{dom} f \cup \operatorname{rng} f)$  such that  $f^n = p(n)$ and  $p(0) = \operatorname{id}_{\operatorname{dom} f \cup \operatorname{rng} f}$  and for every element k of  $\mathbb{N}$  there exists a function g such that g = p(k) and  $p(k+1) = g \cdot f$ .

In the sequel *m*, *n* are natural numbers. We now state a number of propositions:

- (70)  $f^0 = \operatorname{id}_{\operatorname{dom} f \cup \operatorname{rng} f}$ .
- $(71) \quad f^{n+1} = f^n \cdot f.$
- (72)  $f^1 = f$ .

- $(73) \quad f^{n+1} = f \cdot f^n.$
- (74)  $\operatorname{dom}(f^n) \subseteq \operatorname{dom} f \cup \operatorname{rng} f$  and  $\operatorname{rng}(f^n) \subseteq \operatorname{dom} f \cup \operatorname{rng} f$ .
- (75) If  $n \neq 0$ , then dom $(f^n) \subseteq \text{dom } f$  and  $\text{rng}(f^n) \subseteq \text{rng } f$ .
- (76) If rng  $f \subseteq \text{dom } f$ , then  $\text{dom}(f^n) = \text{dom } f$  and rng $(f^n) \subseteq \text{dom } f$ .
- (77)  $f^n \cdot \operatorname{id}_{\operatorname{dom} f \cup \operatorname{rng} f} = f^n$ .
- (78)  $\operatorname{id}_{\operatorname{dom} f \cup \operatorname{rng} f} \cdot f^n = f^n.$
- $(79) \quad f^n \cdot f^m = f^{n+m}.$
- (80) If  $n \neq 0$ , then  $(f^m)^n = f^{m \cdot n}$ .
- (81) If rng  $f \subseteq \text{dom } f$ , then  $(f^m)^n = f^{m \cdot n}$ .
- $(82) \quad \emptyset^n = \emptyset.$
- $(83) \quad (\mathrm{id}_X)^n = \mathrm{id}_X.$
- (84) If rng f misses dom f, then  $f^2 = \emptyset$ .
- (85) For every function f from X into X holds  $f^n$  is a function from X into X.
- (86) For every function f from X into X holds  $f^0 = id_X$ .
- (87) For every function f from X into X holds  $(f^m)^n = f^{m \cdot n}$ .
- (88) For every partial function f from X to X holds  $f^n$  is a partial function from X to X.
- (89) If  $n \neq 0$  and  $a \in X$  and  $f = X \mapsto a$ , then  $f^n = f$ .
- (90) For every function f and for every natural number n holds  $f^n = \text{compose}_{\text{dom } f \cup \text{rng } f}(n \mapsto f)$ .

#### REFERENCES

- [1] Grzegorz Bancerek. Cardinal numbers. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Vol1/card\_1.html.
- [2] Grzegorz Bancerek. The fundamental properties of natural numbers. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Voll/nat\_1.html.
- [3] Grzegorz Bancerek. Curried and uncurried functions. Journal of Formalized Mathematics, 2, 1990. http://mizar.org/JFM/Vol2/ funct\_5.html.
- [4] Grzegorz Bancerek. Reduction relations. Journal of Formalized Mathematics, 7, 1995. http://mizar.org/JFM/Vol7/rewrite1. html.
- [5] Grzegorz Bancerek and Krzysztof Hryniewiecki. Segments of natural numbers and finite sequences. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Vol1/finseq\_1.html.
- [6] Czesław Byliński. Binary operations. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Voll/binop\_1.html.
- [7] Czesław Byliński. Functions and their basic properties. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Voll/funct\_1.html.
- [8] Czesław Byliński. Functions from a set to a set. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Vol1/funct\_ 2.html.
- [9] Czesław Byliński. Partial functions. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Vol1/partfun1.html.
- [10] Czesław Byliński. Some basic properties of sets. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Voll/ zfmisc\_1.html.
- [11] Czesław Byliński. A classical first order language. Journal of Formalized Mathematics, 2, 1990. http://mizar.org/JFM/Vol2/cqc\_lang.html.
- [12] Czesław Byliński. Finite sequences and tuples of elements of a non-empty sets. Journal of Formalized Mathematics, 2, 1990. http: //mizar.org/JFM/Vol2/finseq\_2.html.
- [13] Czesław Byliński. The modification of a function by a function and the iteration of the composition of a function. Journal of Formalized Mathematics, 2, 1990. http://mizar.org/JFM/Vol2/funct\_4.html.

- [14] Czesław Byliński. Cartesian categories. Journal of Formalized Mathematics, 4, 1992. http://mizar.org/JFM/Vol4/cat\_4.html.
- [15] Agata Darmochwał. Finite sets. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Voll/finset\_1.html.
- [16] Beata Madras. Product of family of universal algebras. Journal of Formalized Mathematics, 5, 1993. http://mizar.org/JFM/Vol5/ pralg\_l.html.
- [17] Beata Padlewska. Families of sets. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Voll/setfam\_1.html.
- [18] Dariusz Surowik. Cyclic groups and some of their properties part I. Journal of Formalized Mathematics, 3, 1991. http://mizar. org/JFM/Vol3/gr\_cy\_l.html.
- [19] Andrzej Trybulec. Binary operations applied to functions. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/ Voll/funcop\_1.html.
- [20] Andrzej Trybulec. Domains and their Cartesian products. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/ Voll/domain\_1.html.
- [21] Andrzej Trybulec. Tarski Grothendieck set theory. Journal of Formalized Mathematics, Axiomatics, 1989. http://mizar.org/JFM/ Axiomatics/tarski.html.
- [22] Andrzej Trybulec. Many-sorted sets. Journal of Formalized Mathematics, 5, 1993. http://mizar.org/JFM/Vol5/pboole.html.
- [23] Andrzej Trybulec. Subsets of real numbers. Journal of Formalized Mathematics, Addenda, 2003. http://mizar.org/JFM/Addenda/ numbers.html.
- [24] Wojciech A. Trybulec. Pigeon hole principle. Journal of Formalized Mathematics, 2, 1990. http://mizar.org/JFM/Vol2/finseq\_ 4.html.
- [25] Zinaida Trybulec. Properties of subsets. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Voll/subset\_1.html.
- [26] Edmund Woronowicz. Relations and their basic properties. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/ Voll/relat\_1.html.
- [27] Edmund Woronowicz. Relations defined on sets. Journal of Formalized Mathematics, 1, 1989. http://mizar.org/JFM/Voll/ relset\_1.html.

Received January 12, 1996

Published January 2, 2004