

Function Domains and Frænkel Operator

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Summary. We deal with a non-empty set of functions and a non-empty set of functions from a set A to a non-empty set B . In the case when B is a non-empty set, B^A is redefined. It yields a non-empty set of functions from A to B . An element of such a set is redefined as a function from A to B . Some theorems concerning these concepts are proved, as well as a number of schemes dealing with infinity and the Axiom of Choice. The article contains a number of schemes allowing for simple logical transformations related to terms constructed with the Frænkel Operator.

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The articles [6], [3], [8], [9], [4], [7], [1], [2], and [5] provide the notation and terminology for this paper.

In this paper B is a non empty set and A, C, X are sets.

In this article we present several logical schemes. The scheme *Fraenkel5'* deals with a non empty set \mathcal{A} , a unary functor \mathcal{F} yielding a set, and two unary predicates \mathcal{P} , \mathcal{Q} , and states that:

$$\{\mathcal{F}(v'); v' \text{ ranges over elements of } \mathcal{A} : \mathcal{P}[v']\} \subseteq \{\mathcal{F}(u'); u' \text{ ranges over elements of } \mathcal{A} : \mathcal{Q}[u']\}$$

provided the following condition is satisfied:

- For every element v of \mathcal{A} such that $\mathcal{P}[v]$ holds $\mathcal{Q}[v]$.

The scheme *Fraenkel5''* deals with a non empty set \mathcal{A} , a non empty set \mathcal{B} , a binary functor \mathcal{F} yielding a set, and two binary predicates \mathcal{P} , \mathcal{Q} , and states that:

$$\begin{aligned} &\{\mathcal{F}(u_1, v_1); u_1 \text{ ranges over elements of } \mathcal{A}, v_1 \text{ ranges over elements of } \mathcal{B} : \mathcal{P}[u_1, v_1]\} \subseteq \\ &\{\mathcal{F}(u_2, v_2); u_2 \text{ ranges over elements of } \mathcal{A}, v_2 \text{ ranges over elements of } \mathcal{B} : \mathcal{Q}[u_2, v_2]\} \end{aligned}$$

provided the parameters meet the following requirement:

- For every element u of \mathcal{A} and for every element v of \mathcal{B} such that $\mathcal{P}[u, v]$ holds $\mathcal{Q}[u, v]$.

The scheme *Fraenkel6'* deals with a non empty set \mathcal{A} , a unary functor \mathcal{F} yielding a set, and two unary predicates \mathcal{P} , \mathcal{Q} , and states that:

$$\{\mathcal{F}(v_1); v_1 \text{ ranges over elements of } \mathcal{A} : \mathcal{P}[v_1]\} = \{\mathcal{F}(v_2); v_2 \text{ ranges over elements of } \mathcal{A} : \mathcal{Q}[v_2]\}$$

provided the following condition is satisfied:

- For every element v of \mathcal{A} holds $\mathcal{P}[v]$ iff $\mathcal{Q}[v]$.

The scheme *Fraenkel6''* deals with a non empty set \mathcal{A} , a non empty set \mathcal{B} , a binary functor \mathcal{F} yielding a set, and two binary predicates \mathcal{P} , \mathcal{Q} , and states that:

$$\begin{aligned} &\{\mathcal{F}(u_1, v_1); u_1 \text{ ranges over elements of } \mathcal{A}, v_1 \text{ ranges over elements of } \mathcal{B} : \mathcal{P}[u_1, v_1]\} = \\ &\{\mathcal{F}(u_2, v_2); u_2 \text{ ranges over elements of } \mathcal{A}, v_2 \text{ ranges over elements of } \mathcal{B} : \mathcal{Q}[u_2, v_2]\} \end{aligned}$$

provided the following condition is met:

- For every element u of \mathcal{A} and for every element v of \mathcal{B} holds $\mathcal{P}[u, v]$ iff $\mathcal{Q}[u, v]$.

The scheme *FraenkelF'* deals with a non empty set \mathcal{A} , a unary functor \mathcal{F} yielding a set, a unary functor \mathcal{G} yielding a set, and a unary predicate \mathcal{P} , and states that:

$$\{\mathcal{F}(v_1); v_1 \text{ ranges over elements of } \mathcal{A} : \mathcal{P}[v_1]\} = \{\mathcal{G}(v_2); v_2 \text{ ranges over elements of } \mathcal{A} : \mathcal{P}[v_2]\}$$

provided the parameters have the following property:

- For every element v of \mathcal{A} holds $\mathcal{F}(v) = \mathcal{G}(v)$.

The scheme *FraenkelF'R* deals with a non empty set \mathcal{A} , a unary functor \mathcal{F} yielding a set, a unary functor \mathcal{G} yielding a set, and a unary predicate \mathcal{P} , and states that:

$$\{\mathcal{F}(v_1); v_1 \text{ ranges over elements of } \mathcal{A} : \mathcal{P}[v_1]\} = \{\mathcal{G}(v_2); v_2 \text{ ranges over elements of } \mathcal{A} : \mathcal{P}[v_2]\}$$

provided the following condition is met:

- For every element v of \mathcal{A} such that $\mathcal{P}[v]$ holds $\mathcal{F}(v) = \mathcal{G}(v)$.

The scheme *FraenkelF''* deals with a non empty set \mathcal{A} , a non empty set \mathcal{B} , a binary functor \mathcal{F} yielding a set, a binary functor \mathcal{G} yielding a set, and a binary predicate \mathcal{P} , and states that:

$$\{\mathcal{F}(u_1, v_1); u_1 \text{ ranges over elements of } \mathcal{A}, v_1 \text{ ranges over elements of } \mathcal{B} : \mathcal{P}[u_1, v_1]\} = \{\mathcal{G}(u_2, v_2); u_2 \text{ ranges over elements of } \mathcal{A}, v_2 \text{ ranges over elements of } \mathcal{B} : \mathcal{P}[u_2, v_2]\}$$

provided the following condition is satisfied:

- For every element u of \mathcal{A} and for every element v of \mathcal{B} holds $\mathcal{F}(u, v) = \mathcal{G}(u, v)$.

The scheme *FraenkelF6''* deals with a non empty set \mathcal{A} , a non empty set \mathcal{B} , a binary functor \mathcal{F} yielding a set, and two binary predicates \mathcal{P} , \mathcal{Q} , and states that:

$$\{\mathcal{F}(u_1, v_1); u_1 \text{ ranges over elements of } \mathcal{A}, v_1 \text{ ranges over elements of } \mathcal{B} : \mathcal{P}[u_1, v_1]\} = \{\mathcal{F}(v_2, u_2); u_2 \text{ ranges over elements of } \mathcal{A}, v_2 \text{ ranges over elements of } \mathcal{B} : \mathcal{Q}[u_2, v_2]\}$$

provided the parameters have the following properties:

- For every element u of \mathcal{A} and for every element v of \mathcal{B} holds $\mathcal{P}[u, v]$ iff $\mathcal{Q}[u, v]$, and
- For every element u of \mathcal{A} and for every element v of \mathcal{B} holds $\mathcal{F}(u, v) = \mathcal{F}(v, u)$.

The following propositions are true:

(3)¹ Let A, B be non empty sets, F, G be functions from A into B , and X be a set. If $F \upharpoonright X = G \upharpoonright X$, then for every element x of A such that $x \in X$ holds $F(x) = G(x)$.

(5)² For all sets A, B holds $B^A \subseteq 2^{[A, B]}$.

(6) For all sets X, Y such that $Y^X \neq \emptyset$ and $X \subseteq A$ and $Y \subseteq B$ holds every element of Y^X is a partial function from A to B .

Now we present a number of schemes. The scheme *RelevantArgs* deals with a non empty set \mathcal{A} , a non empty set \mathcal{B} , a set \mathcal{C} , a function \mathcal{D} from \mathcal{A} into \mathcal{B} , a function \mathcal{E} from \mathcal{A} into \mathcal{B} , and two unary predicates \mathcal{P} , \mathcal{Q} , and states that:

$$\{\mathcal{D}(u'); u' \text{ ranges over elements of } \mathcal{A} : \mathcal{P}[u'] \wedge u' \in \mathcal{C}\} = \{\mathcal{E}(v'); v' \text{ ranges over elements of } \mathcal{A} : \mathcal{Q}[v'] \wedge v' \in \mathcal{C}\}$$

provided the parameters meet the following requirements:

- $\mathcal{D} \upharpoonright \mathcal{C} = \mathcal{E} \upharpoonright \mathcal{C}$, and
- For every element u of \mathcal{A} such that $u \in \mathcal{C}$ holds $\mathcal{P}[u]$ iff $\mathcal{Q}[u]$.

The scheme *Fr Set0* deals with a non empty set \mathcal{A} and a unary predicate \mathcal{P} , and states that:

$$\{x; x \text{ ranges over elements of } \mathcal{A} : \mathcal{P}[x]\} \subseteq \mathcal{A}$$

for all values of the parameters.

The scheme *GenI''* deals with a non empty set \mathcal{A} , a non empty set \mathcal{B} , a binary functor \mathcal{F} yielding a set, a unary predicate \mathcal{Q} , and a binary predicate \mathcal{P} , and states that:

$$\text{For every element } s \text{ of } \mathcal{A} \text{ and for every element } t \text{ of } \mathcal{B} \text{ such that } \mathcal{P}[s, t] \text{ holds } \mathcal{Q}[\mathcal{F}(s, t)]$$

provided the parameters satisfy the following condition:

- For every set s_1 such that $s_1 \in \{\mathcal{F}(s_2, t_1); s_2 \text{ ranges over elements of } \mathcal{A}, t_1 \text{ ranges over elements of } \mathcal{B} : \mathcal{P}[s_2, t_1]\}$ holds $\mathcal{Q}[s_1]$.

The scheme *GenI''A* deals with a non empty set \mathcal{A} , a non empty set \mathcal{B} , a binary functor \mathcal{F} yielding a set, a unary predicate \mathcal{Q} , and a binary predicate \mathcal{P} , and states that:

¹ The propositions (1) and (2) have been removed.

² The proposition (4) has been removed.

For every set s_1 such that $s_1 \in \{\mathcal{F}(s_2, t_1); s_2 \text{ ranges over elements of } \mathcal{A}, t_1 \text{ ranges over elements of } \mathcal{B} : \mathcal{P}[s_2, t_1]\}$ holds $Q[s_1]$

provided the following condition is met:

- For every element s of \mathcal{A} and for every element t of \mathcal{B} such that $\mathcal{P}[s, t]$ holds $Q[\mathcal{F}(s, t)]$.

The scheme *Gen2*'' deals with a non empty set \mathcal{A} , a non empty set \mathcal{B} , a non empty set \mathcal{C} , a binary functor \mathcal{F} yielding an element of \mathcal{C} , a unary predicate Q , and a binary predicate \mathcal{P} , and states that:

$$\{s_1; s_1 \text{ ranges over elements of } \mathcal{C} : s_1 \in \{\mathcal{F}(s_2, t_1); s_2 \text{ ranges over elements of } \mathcal{A}, t_1 \text{ ranges over elements of } \mathcal{B} : \mathcal{P}[s_2, t_1]\} \wedge Q[s_1]\} = \{\mathcal{F}(s_3, t_2); s_3 \text{ ranges over elements of } \mathcal{A}, t_2 \text{ ranges over elements of } \mathcal{B} : \mathcal{P}[s_3, t_2] \wedge Q[\mathcal{F}(s_3, t_2)]\}$$

for all values of the parameters.

The scheme *Gen3*' deals with a non empty set \mathcal{A} , a unary functor \mathcal{F} yielding a set, and two unary predicates \mathcal{P} , Q , and states that:

$$\{\mathcal{F}(s); s \text{ ranges over elements of } \mathcal{A} : s \in \{s_2; s_2 \text{ ranges over elements of } \mathcal{A} : Q[s_2]\} \wedge \mathcal{P}[s]\} = \{\mathcal{F}(s_3); s_3 \text{ ranges over elements of } \mathcal{A} : Q[s_3] \wedge \mathcal{P}[s_3]\}$$

for all values of the parameters.

The scheme *Gen3*'' deals with a non empty set \mathcal{A} , a non empty set \mathcal{B} , a binary functor \mathcal{F} yielding a set, a unary predicate Q , and a binary predicate \mathcal{P} , and states that:

$$\{\mathcal{F}(s, t); s \text{ ranges over elements of } \mathcal{A}, t \text{ ranges over elements of } \mathcal{B} : s \in \{s_2; s_2 \text{ ranges over elements of } \mathcal{A} : Q[s_2]\} \wedge \mathcal{P}[s, t]\} = \{\mathcal{F}(s_3, t_2); s_3 \text{ ranges over elements of } \mathcal{A}, t_2 \text{ ranges over elements of } \mathcal{B} : Q[s_3] \wedge \mathcal{P}[s_3, t_2]\}$$

for all values of the parameters.

The scheme *Gen4*'' deals with a non empty set \mathcal{A} , a non empty set \mathcal{B} , a binary functor \mathcal{F} yielding a set, and two binary predicates \mathcal{P} , Q , and states that:

$$\{\mathcal{F}(s, t); s \text{ ranges over elements of } \mathcal{A}, t \text{ ranges over elements of } \mathcal{B} : \mathcal{P}[s, t]\} \subseteq \{\mathcal{F}(s_2, t_1); s_2 \text{ ranges over elements of } \mathcal{A}, t_1 \text{ ranges over elements of } \mathcal{B} : Q[s_2, t_1]\}$$

provided the parameters meet the following condition:

- Let s be an element of \mathcal{A} and t be an element of \mathcal{B} . If $\mathcal{P}[s, t]$, then there exists an element s' of \mathcal{A} such that $Q[s', t]$ and $\mathcal{F}(s, t) = \mathcal{F}(s', t)$.

The scheme *FrSet 1* deals with a non empty set \mathcal{A} , a set \mathcal{B} , a unary functor \mathcal{F} yielding a set, and a unary predicate \mathcal{P} , and states that:

$$\{\mathcal{F}(y); y \text{ ranges over elements of } \mathcal{A} : \mathcal{F}(y) \in \mathcal{B} \wedge \mathcal{P}[y]\} \subseteq \mathcal{B}$$

for all values of the parameters.

The scheme *FrSet 2* deals with a non empty set \mathcal{A} , a set \mathcal{B} , a unary functor \mathcal{F} yielding a set, and a unary predicate \mathcal{P} , and states that:

$$\{\mathcal{F}(y); y \text{ ranges over elements of } \mathcal{A} : \mathcal{P}[y] \wedge \mathcal{F}(y) \notin \mathcal{B}\} \text{ misses } \mathcal{B}$$

for all values of the parameters.

The scheme *FrEqual 1* deals with a non empty set \mathcal{A} , a non empty set \mathcal{B} , a binary functor \mathcal{F} yielding a set, an element C of \mathcal{B} , and two binary predicates \mathcal{P} , Q , and states that:

$$\{\mathcal{F}(s, t); s \text{ ranges over elements of } \mathcal{A}, t \text{ ranges over elements of } \mathcal{B} : Q[s, t]\} = \{\mathcal{F}(s', C); s' \text{ ranges over elements of } \mathcal{A} : \mathcal{P}[s', C]\}$$

provided the following condition is satisfied:

- For every element s of \mathcal{A} and for every element t of \mathcal{B} holds $Q[s, t]$ iff $t = C$ and $\mathcal{P}[s, t]$.

The scheme *FrEqual 2* deals with a non empty set \mathcal{A} , a non empty set \mathcal{B} , a binary functor \mathcal{F} yielding a set, an element C of \mathcal{B} , and a binary predicate \mathcal{P} , and states that:

$$\{\mathcal{F}(s, t); s \text{ ranges over elements of } \mathcal{A}, t \text{ ranges over elements of } \mathcal{B} : t = C \wedge \mathcal{P}[s, t]\} = \{\mathcal{F}(s', C); s' \text{ ranges over elements of } \mathcal{A} : \mathcal{P}[s', C]\}$$

for all values of the parameters.

Let I_1 be a set. We say that I_1 is functional if and only if:

(Def. 1) For every set x such that $x \in I_1$ holds x is a function.

One can verify that there exists a set which is non empty and functional.

Let P be a functional set. Observe that every element of P is function-like and relation-like.

The following proposition is true

(8)³ For every function f holds $\{f\}$ is functional.

Let A, B be sets. Observe that B^A is functional.

Let A, B be sets. A functional non empty set is said to be a non empty set of functions from A to B if:

(Def. 2) Every element of it is a function from A into B .

We now state two propositions:

(10)⁴ For every function f from A into C holds $\{f\}$ is a non empty set of functions from A to C .

(11) B^A is a non empty set of functions from A to B .

Let A be a set and let B be a non empty set. Then B^A is a non empty set of functions from A to B . Let F be a non empty set of functions from A to B . We see that the element of F is a function from A into B .

In the sequel p_1 is an element of B^A .

We now state two propositions:

(14)⁵ Let X, Y be sets. Suppose $Y^X \neq \emptyset$ and $X \subseteq A$ and $Y \subseteq B$. Let f be an element of Y^X . Then there exists an element p_1 of B^A such that $p_1 \upharpoonright X = f$.

(15) For every set X and for every p_1 holds $p_1 \upharpoonright X = p_1 \upharpoonright (A \cap X)$.

Now we present four schemes. The scheme *FraenkelFin* deals with a non empty set \mathcal{A} , a set \mathcal{B} , and a unary functor \mathcal{F} yielding a set, and states that:

$\{\mathcal{F}(w); w \text{ ranges over elements of } \mathcal{A} : w \in \mathcal{B}\}$ is finite

provided the parameters meet the following requirement:

- \mathcal{B} is finite.

The scheme *CartFin* deals with non empty sets \mathcal{A}, \mathcal{B} , sets \mathcal{C}, \mathcal{D} , and a binary functor \mathcal{F} yielding a set, and states that:

$\{\mathcal{F}(u', v'); u' \text{ ranges over elements of } \mathcal{A}, v' \text{ ranges over elements of } \mathcal{B} : u' \in \mathcal{C} \wedge v' \in \mathcal{D}\}$ is finite

provided the parameters meet the following requirements:

- \mathcal{C} is finite, and
- \mathcal{D} is finite.

The scheme *Finiteness* deals with a non empty set \mathcal{A} , an element \mathcal{B} of $\text{Fin } \mathcal{A}$, and a binary predicate \mathcal{P} , and states that:

Let x be an element of \mathcal{A} . Suppose $x \in \mathcal{B}$. Then there exists an element y of \mathcal{A} such that $y \in \mathcal{B}$ and $\mathcal{P}[y, x]$ and for every element z of \mathcal{A} such that $z \in \mathcal{B}$ and $\mathcal{P}[z, y]$ holds $\mathcal{P}[y, z]$

provided the following conditions are satisfied:

- For every element x of \mathcal{A} holds $\mathcal{P}[x, x]$, and
- For all elements x, y, z of \mathcal{A} such that $\mathcal{P}[x, y]$ and $\mathcal{P}[y, z]$ holds $\mathcal{P}[x, z]$.

The scheme *Fin Im* deals with a non empty set \mathcal{A} , a non empty set \mathcal{B} , an element C of $\text{Fin } \mathcal{B}$, a unary functor \mathcal{F} yielding an element of \mathcal{A} , and a binary predicate \mathcal{P} , and states that:

There exists an element c_1 of $\text{Fin } \mathcal{A}$ such that for every element t of \mathcal{A} holds $t \in c_1$ if and only if there exists an element t' of \mathcal{B} such that $t' \in C$ and $t = \mathcal{F}(t')$ and $\mathcal{P}[t, t']$

for all values of the parameters.

The following proposition is true

(16) For all sets A, B such that A is finite and B is finite holds B^A is finite.

³ The proposition (7) has been removed.

⁴ The proposition (9) has been removed.

⁵ The propositions (12) and (13) have been removed.

Now we present three schemes. The scheme *ImFin* deals with a non empty set \mathcal{A} , a non empty set \mathcal{B} , a set \mathcal{C} , a set \mathcal{D} , and a unary functor \mathcal{F} yielding a set, and states that:

$\{\mathcal{F}(p'_1); p'_1 \text{ ranges over elements of } \mathcal{B}^{\mathcal{A}}: p'_1 \circ \mathcal{C} \subseteq \mathcal{D}\}$ is finite

provided the parameters meet the following conditions:

- \mathcal{C} is finite,
- \mathcal{D} is finite, and
- For all elements p_1, p_2 of $\mathcal{B}^{\mathcal{A}}$ such that $p_1 \upharpoonright \mathcal{C} = p_2 \upharpoonright \mathcal{C}$ holds $\mathcal{F}(p_1) = \mathcal{F}(p_2)$.

The scheme *FunctChoice* deals with a non empty set \mathcal{A} , a non empty set \mathcal{B} , an element C of $\text{Fin } \mathcal{A}$, and a binary predicate \mathcal{P} , and states that:

There exists a function f_1 from \mathcal{A} into \mathcal{B} such that for every element t of \mathcal{A} such that $t \in C$ holds $\mathcal{P}[t, f_1(t)]$

provided the following requirement is met:

- For every element t of \mathcal{A} such that $t \in C$ there exists an element f_1 of \mathcal{B} such that $\mathcal{P}[t, f_1]$.

The scheme *FuncsChoice* deals with a non empty set \mathcal{A} , a non empty set \mathcal{B} , an element C of $\text{Fin } \mathcal{A}$, and a binary predicate \mathcal{P} , and states that:

There exists an element f_1 of $\mathcal{B}^{\mathcal{A}}$ such that for every element t of \mathcal{A} such that $t \in C$ holds $\mathcal{P}[t, f_1(t)]$

provided the parameters meet the following requirement:

- For every element t of \mathcal{A} such that $t \in C$ there exists an element f_1 of \mathcal{B} such that $\mathcal{P}[t, f_1]$.

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