## **Definitions of Petri Net. Part I**

## Waldemar Korczyński Pedagogical University Kielce

**Summary.** In the paper the classical definition of Petri net is described. The article also contains some theorems needed for proving equivalences of these definitions with other definitions of Petri net as relational algebras. See [3], [4] and other.

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The articles [5], [1], [6], and [2] provide the notation and terminology for this paper. In this paper *x*, *y*, *X*, *Y* denote sets. Let *N* be a net. The functor chaos(*N*) yields a set and is defined by:

 $(\text{Def. 2})^1$  chaos $(N) = \text{Elements}(N) \cup \{\text{Elements}(N)\}.$ 

In the sequel *M* is a Petri net.

Let us consider X, Y. Let us assume that X misses Y. The functor  $\operatorname{PTempty}_f(X,Y)$  yielding a strict Petri net is defined by:

(Def. 4)<sup>2</sup> PTempty<sub>f</sub>(X,Y) =  $\langle X, Y, \emptyset \rangle$ .

Let us consider X. The functor Tempty f(X) yields a strict Petri net and is defined by:

(Def. 5) Tempty<sub>f</sub>(X) = PTempty<sub>f</sub>(X, $\emptyset$ ).

The functor  $\operatorname{Pempty}_{f}(X)$  yields a strict Petri net and is defined by:

(Def. 6) Pempty<sub>*f*</sub>(X) = PTempty<sub>*f*</sub>( $\emptyset$ ,X).

Let us consider x. The functor  $Tsingle_f(x)$  yields a strict Petri net and is defined as follows:

(Def. 7) Tsingle<sub>f</sub>(x) = PTempty<sub>f</sub>( $\emptyset$ , {x}).

The functor  $Psingle_f(x)$  yields a strict Petri net and is defined by:

(Def. 8) Psingle  $_{f}(x) = \text{PTempty}_{f}(\{x\}, \emptyset).$ 

The strict Petri net empty  $_{f}$  is defined as follows:

(Def. 9) empty  $_f = \text{PTempty}_f(\emptyset, \emptyset)$ .

One can prove the following propositions:

<sup>&</sup>lt;sup>1</sup> The definition (Def. 1) has been removed.

<sup>&</sup>lt;sup>2</sup> The definition (Def. 3) has been removed.

- (2)<sup>3</sup> Suppose X misses Y. Then the places of  $\operatorname{PTempty}_f(X,Y) = X$  and the transitions of  $\operatorname{PTempty}_f(X,Y) = Y$  and the flow relation of  $\operatorname{PTempty}_f(X,Y) = \emptyset$ .
- (3) The places of Tempty<sub>f</sub>(X) = X and the transitions of Tempty<sub>f</sub>(X) =  $\emptyset$  and the flow relation of Tempty<sub>f</sub>(X) =  $\emptyset$ .
- (4) For every X holds the places of  $\text{Pempty}_f(X) = \emptyset$  and the transitions of  $\text{Pempty}_f(X) = X$  and the flow relation of  $\text{Pempty}_f(X) = \emptyset$ .
- (5) For every x holds the places of  $\text{Tsingle}_f(x) = \emptyset$  and the transitions of  $\text{Tsingle}_f(x) = \{x\}$  and the flow relation of  $\text{Tsingle}_f(x) = \emptyset$ .
- (6) For every x holds the places of Psingle<sub>f</sub>(x) = {x} and the transitions of Psingle<sub>f</sub>(x) = Ø and the flow relation of Psingle<sub>f</sub>(x) = Ø.
- (7) The places of  $empty_f = \emptyset$  and the transitions of  $empty_f = \emptyset$  and the flow relation of  $empty_f = \emptyset$ .
- (8) The places of  $M \subseteq \text{Elements}(M)$  and the transitions of  $M \subseteq \text{Elements}(M)$ .
- (11)<sup>4</sup>(i) If  $\langle x, y \rangle \in$  the flow relation of M and  $x \in$  the transitions of M, then  $x \notin$  the places of M and  $y \notin$  the transitions of M and  $y \in$  the places of M,
- (ii) if  $\langle x, y \rangle \in$  the flow relation of *M* and  $y \in$  the transitions of *M*, then  $y \notin$  the places of *M* and  $x \notin$  the transitions of *M* and  $x \in$  the places of *M*,
- (iii) if  $\langle x, y \rangle \in$  the flow relation of *M* and  $x \in$  the places of *M*, then  $y \notin$  the places of *M* and  $x \notin$  the transitions of *M* and  $y \in$  the transitions of *M*, and
- (iv) if  $\langle x, y \rangle \in$  the flow relation of *M* and  $y \in$  the places of *M*, then  $x \notin$  the places of *M* and  $y \notin$  the transitions of *M* and  $x \in$  the transitions of *M*.
- (12) chaos(M)  $\neq \emptyset$ .
- (13) The flow relation of  $M \subseteq [: \text{Elements}(M), \text{Elements}(M):]$  and (the flow relation of  $M)^{\smile} \subseteq [: \text{Elements}(M), \text{Elements}(M):].$
- (14)  $\operatorname{rng}((\operatorname{the flow relation of } M) \upharpoonright (\operatorname{the transitions of } M)) \subseteq \operatorname{the places of } M \operatorname{and rng}((\operatorname{the flow relation of } M) \cong \operatorname{the transitions of } M) \subseteq \operatorname{the places of } M \operatorname{and dom}((\operatorname{the flow relation of } M) \upharpoonright (\operatorname{the transitions of } M)) \subseteq \operatorname{the transitions of } M \operatorname{and rng}((\operatorname{the flow relation of } M) \cong \operatorname{the transitions of } M \operatorname{and rng}((\operatorname{the flow relation of } M)) \subseteq \operatorname{the transitions of } M \operatorname{and rng}((\operatorname{the flow relation of } M)) \subseteq \operatorname{the transitions of } M \operatorname{and rng}((\operatorname{the flow relation of } M)) \subseteq \operatorname{the transitions of } M \operatorname{and rng}((\operatorname{the flow relation of } M)) \subseteq \operatorname{the transitions of } M \operatorname{and rng}((\operatorname{the flow relation of } M)) \subseteq \operatorname{the transitions of } M \operatorname{and rng}((\operatorname{the flow relation of } M)) \subseteq \operatorname{the transitions of } M \operatorname{and rng}((\operatorname{the flow relation of } M)) \subseteq \operatorname{the transitions of } M \operatorname{and rng}((\operatorname{the flow relation of } M)) \subseteq \operatorname{the transitions of } M \operatorname{and rng}(\operatorname{the flow relation of } M) \subseteq \operatorname{the transitions of } M \operatorname{and rng}(\operatorname{the flow relation of } M) \subseteq \operatorname{the places of } M \operatorname{and rng}(\operatorname{the flow relation of } M) \subseteq \operatorname{the transitions of } M \operatorname{and rng}(\operatorname{the flow relation of } M) \subseteq \operatorname{the transitions of } M \operatorname{and rng}(\operatorname{the flow relation of } M) \subseteq \operatorname{the places of } M \operatorname{and rng}(\operatorname{the flow relation of } M) \subseteq \operatorname{the transitions of } M \operatorname{and rng}(\operatorname{the flow relation of } M) \subseteq \operatorname{the places of } M \operatorname{and rng}(\operatorname{the flow relation of } M) \subseteq \operatorname{the places of } M \operatorname{and rng}(\operatorname{the flow relation of } M) \subseteq \operatorname{the places of } M \operatorname{and rng}(\operatorname{the flow relation of } M) \subseteq \operatorname{the places of } M \operatorname{and rng}(\operatorname{the flow relation of } M) \subseteq \operatorname{the places of } M \operatorname{and rng}(\operatorname{the flow relation of } M) \subseteq \operatorname{the places of } M \operatorname{and rng}(\operatorname{the flow relation of } M) \subseteq \operatorname{the places of } M \operatorname{and rng}(\operatorname{the flow relation of } M) \subseteq \operatorname{the places of } M \operatorname{and rng}(\operatorname{the flow relation } M) \subseteq \operatorname{the places of } M \operatorname{and rng}(\operatorname{the flow relation } M) \subseteq \operatorname{the places of } M \operatorname{and rng}(\operatorname{the flow relation } M) \subseteq \operatorname{the places of } M \operatorname{and rng}(\operatorname{the flow relation } M) \subseteq \operatorname{the places of } M \operatorname{and rng}(\operatorname{th flow relation } M) \subseteq \operatorname{the plac$
- (15) rng((the flow relation of *M*) ↾(the transitions of *M*)) misses dom((the flow relation of *M*) ↾(the transitions of *M*)) and rng((the flow relation of *M*) ↾(the transitions of *M*)) misses dom((the flow relation of *M*) ∼ ↾the transitions of *M*) and rng((the flow relation of *M*) ↾(the transitions of *M*)) misses dom((the flow relation of *M*) ∼ ↾the transitions of *M*) and rng((the flow relation of *M*) ∼ ↾the transitions of *M*) and rng((the flow relation of *M*) ∼ ↾the transitions of *M*)) and rng((the flow relation of *M*) ∼ ↾the transitions of *M*) misses dom((the flow relation of *M*) ↾(the transitions of *M*)) and rng((the flow relation of *M*) ∼ ↾the transitions of *M*) misses dom((the flow relation of *M*) ∼ ↾the transitions of *M*) misses dom((the flow relation of *M*) ∼ ↾the transitions of *M*) misses dom((the flow relation of *M*) ∼ ↾the transitions of *M*) misses dom((the flow relation of *M*) ∼ ↾the transitions of *M*) misses dom((the flow relation of *M*) ∼ ↾the transitions of *M*) misses dom((the flow relation of *M*) ↾ (the transitions of *M*)) misses rng((the flow relation of *M*) ↾ (the transitions of *M*)) misses rng((the flow relation of *M*) ∼ ↾the transitions of *M*) and dom((the flow relation of *M*) ↾ (the transitions of *M*)) misses rng((the flow relation of *M*) ∼ ↾ the transitions of *M*) and dom((the flow relation of *M*) ↾ (the transitions of *M*)) misses rng((the flow relation of *M*) ⊨ (the transitions of *M*)) misses rng((the flow relation of *M*) ↾ (the transitions of *M*)) misses rng((the flow relation of *M*) ⊨ (the transitions of *M*)) misses rng((the flow relation of *M*) ⊨ (the transitions of *M*)) misses rng((the flow relation of *M*) ⊨ (the transitions of *M*)) misses rng((the flow relation of *M*) ⊨ (the transitions of *M*)) misses rng((the flow relation of *M*) ⊨ (the transitions of *M*)) misses rng((the flow relation of *M*) ⊨ (the transitions of *M*)) misses rng((the flow relation of *M*) ⊨ (the transitions of *M*)) misses rng((the flow relation of *M*)

<sup>&</sup>lt;sup>3</sup> The proposition (1) has been removed.

<sup>&</sup>lt;sup>4</sup> The propositions (9) and (10) have been removed.

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- (16) ((The flow relation of *M*) \(the transitions of *M*)) · ((the flow relation of *M*) \(the transitions of *M*)) = Ø and ((the flow relation of *M*) \(the transitions of *M*)) · ((the flow relation of *M*)) \(the transitions of *M*)) · ((the flow relation of *M*)) \(the transitions of *M*)) · ((the flow relation of *M*)) \(the transitions of *M*)) · ((the flow relation of *M*)) \(the transitions of *M*)) = Ø and ((the flow relation of *M*)) \(the transitions of *M*)) · ((the flow relation of *M*)) \(the transitions of *M*)) = Ø and ((the flow relation of *M*)) \(the places of *M*)) · ((the flow relation of *M*)) \(the places of *M*)) = Ø and ((the flow relation of *M*)) \(the places of *M*)) · ((the flow relation of *M*)) \(the places of *M*)) = Ø and ((the flow relation of *M*)) \(the places of *M*)) · ((the flow relation of *M*)) \(the places of *M*)) = Ø and ((the flow relation of *M*)) \(the places of *M*)) · ((the flow relation of *M*)) \(the places of *M*)) = Ø and ((the flow relation of *M*)) \(the places of *M*)) = Ø and ((the flow relation of *M*)) \(the places of *M*)) · ((the flow relation of *M*)) \(the places of *M*) = Ø and ((the flow relation of *M*)) \(the places of *M*)) · ((the flow relation of *M*)) \(the places of *M*) = Ø and ((the flow relation of *M*)) \(the places of *M*)) · ((the flow relation of *M*)) \(the places of *M*) = Ø and ((the flow relation of *M*)) \(the places of *M*)) · ((the flow relation of *M*)) \(the places of *M*) = Ø.
- (17) ((The flow relation of M)  $\uparrow$  (the transitions of M))  $\cdot$  id<sub>the places of M</sub> = (the flow relation of M) (the transitions of M) and (the flow relation of M)  $\subset$  (the transitions of M).  $\operatorname{id_{the places of M}} = (\text{the flow relation of } M)^{\smile} \upharpoonright \text{the transitions of } M \text{ and } \operatorname{id_{the transitions of } M} \cdot ((\text{the transitions of } M)^{\smile})^{\smile} \upharpoonright M$ flow relation of M) (the transitions of M) = (the flow relation of M) (the transitions of M) and  $\operatorname{id}_{\operatorname{the transitions of } M} \cdot ((\operatorname{the flow relation of } M)) \subset |\operatorname{the transitions of } M) = (\operatorname{the flow relation of } M)$ M)  $\sim$  [the transitions of M and ((the flow relation of M) [(the places of M))  $\cdot$  id<sub>the transitions of M =</sub> (the flow relation of M)  $\restriction$  (the places of M) and ((the flow relation of M)  $\checkmark$   $\restriction$  the places of M)  $\cdot$  id<sub>the transitions of M = (the flow relation of M)  $\subset$  [the places of M and id<sub>the places of M  $\cdot$  ((the</sub></sub> flow relation of M (the places of M) = (the flow relation of M) (the places of M) and  $\operatorname{id_{the places of M}} \cdot ((\text{the flow relation of } M) \subset [\text{the places of } M) = (\text{the flow relation of } M) \subset [\text{the places of } M]$ places of M and ((the flow relation of M)) (the places of M))  $\cdot id_{\text{the transitions of }M} = (the$ flow relation of M) (the places of M) and ((the flow relation of M) (the places of M).  $\mathrm{id}_{\mathrm{the transitions of }M} = (\mathrm{the flow relation of }M) \subset \mathrm{the places of }M \mathrm{ and id}_{\mathrm{the transitions of }M} \cdot ((\mathrm{the flow flow flow }M))$ relation of M)  $\uparrow$  (the places of M)) =  $\emptyset$  and id<sub>the transitions of  $M \cdot ($ (the flow relation of  $M) \subset \uparrow$  the</sub> places of M) =  $\emptyset$  and ((the flow relation of M)  $\uparrow$  (the places of M))  $\cdot$  id<sub>the places of M</sub> =  $\emptyset$  and ((the flow relation of M)  $\subset$  [the places of M)  $\cdot$  id<sub>the places of  $M = \emptyset$  and id<sub>the places of M  $\cdot$  ((the flow relation of M)  $\cap$  ((t</sub></sub> tion of M (the transitions of M) =  $\emptyset$  and id<sub>the places of M</sub> ((the flow relation of M)  $\check{}$  [the transitions of M) =  $\emptyset$  and ((the flow relation of M))(the transitions of M))  $\cdot$  id<sub>the transitions of M =  $\emptyset$ </sub> and ((the flow relation of M)  $\simeq$  [the transitions of M)  $\cdot$  id<sub>the transitions of  $M = \emptyset$ .</sub>
- (18)(i) (The flow relation of M)  $\subset$  [the transitions of M misses  $id_{Elements(M)}$ ,
- (ii) (the flow relation of M)  $\uparrow$  (the transitions of M) misses  $id_{Elements(M)}$ ,
- (iii) (the flow relation of M)  $\subset$  [the places of M misses id<sub>Elements(M)</sub>, and
- (iv) (the flow relation of M)  $\uparrow$  (the places of M) misses  $id_{Elements(M)}$ .

- (20)(i) ((The flow relation of M) $\restriction$ (the places of M)) $\sim$  = (the flow relation of M) $\sim$  $\restriction$ the transitions of M, and
- (ii) ((the flow relation of M)  $\upharpoonright$  (the transitions of M))  $\simeq$  = (the flow relation of M)  $\cong$   $\upharpoonright$  the places of M.
- (21)(i) (The flow relation of M)  $\restriction$  (the places of M)  $\cup$  (the flow relation of M)  $\restriction$  (the transitions of M) = the flow relation of M,
- (ii) (the flow relation of *M*) ↾(the transitions of *M*) ∪ (the flow relation of *M*) ↾(the places of *M*) = the flow relation of *M*,
- (iii) ((the flow relation of M)  $\uparrow$  (the places of M))  $\sim \cup$  ((the flow relation of M)  $\uparrow$  (the transitions of M))  $\sim =$  (the flow relation of M)  $\sim$ , and
- (iv) ((the flow relation of M) $\restriction$ (the transitions of M)) $\sim \cup$  ((the flow relation of M) $\restriction$ (the places of M)) $\sim =$  (the flow relation of M) $\sim$ .

Let us consider *M*. The functor enter $_f(M)$  yields a binary relation and is defined by:

(Def. 10) enter<sub>*f*</sub>(*M*) = (the flow relation of *M*)  $\subset$  [the transitions of *M*  $\cup$  id<sub>the places of *M*.</sub>

The functor  $\operatorname{exit}_f(M)$  yields a binary relation and is defined as follows:

(Def. 11)  $\operatorname{exit}_f(M) = (\text{the flow relation of } M) \upharpoonright (\text{the transitions of } M) \cup \operatorname{id}_{\operatorname{the places of } M}.$ 

The following four propositions are true:

- (22)  $\operatorname{exit}_{f}(M) \subseteq [:\operatorname{Elements}(M), \operatorname{Elements}(M):]$  and  $\operatorname{enter}_{f}(M) \subseteq [:\operatorname{Elements}(M), \operatorname{Elements}(M):]$ .
- (23) domexit<sub>f</sub>(M)  $\subseteq$  Elements(M) and rngexit<sub>f</sub>(M)  $\subseteq$  Elements(M) and domenter<sub>f</sub>(M)  $\subseteq$  Elements(M) and rngenter<sub>f</sub>(M)  $\subseteq$  Elements(M).
- (24)  $\operatorname{exit}_f(M) \cdot \operatorname{exit}_f(M) = \operatorname{exit}_f(M)$  and  $\operatorname{exit}_f(M) \cdot \operatorname{enter}_f(M) = \operatorname{exit}_f(M)$  and  $\operatorname{enter}_f(M) \cdot \operatorname{exit}_f(M) = \operatorname{enter}_f(M)$ .
- (25)  $\operatorname{exit}_{f}(M) \cdot (\operatorname{exit}_{f}(M) \setminus \operatorname{id}_{\operatorname{Elements}(M)}) = \emptyset$  and  $\operatorname{enter}_{f}(M) \cdot (\operatorname{enter}_{f}(M) \setminus \operatorname{id}_{\operatorname{Elements}(M)}) = \emptyset$ .

Let us consider *M*. The functor  $\text{prox}_f(M)$  yielding a binary relation is defined by the condition (Def. 12).

(Def. 12)  $\operatorname{prox}_f(M) = (\text{the flow relation of } M) \upharpoonright (\text{the places of } M) \cup (\text{the flow relation of } M)^{\sim} \upharpoonright \text{the places of } M \cup \operatorname{id}_{\operatorname{the places of } M}.$ 

The functor flow f(M) yielding a binary relation is defined by:

(Def. 13) flow<sub>*f*</sub>(*M*) = (the flow relation of *M*)  $\cup$  id<sub>Elements(*M*)</sub>.

We now state the proposition

(26)  $\operatorname{prox}_f(M) \cdot \operatorname{prox}_f(M) = \operatorname{prox}_f(M)$  and  $(\operatorname{prox}_f(M) \setminus \operatorname{id}_{\operatorname{Elements}(M)}) \cdot \operatorname{prox}_f(M) = \emptyset$  and  $\operatorname{prox}_f(M) \cup (\operatorname{prox}_f(M))^{\smile} \cup \operatorname{id}_{\operatorname{Elements}(M)} = \operatorname{flow}_f(M) \cup (\operatorname{flow}_f(M))^{\smile}$ .

Let us consider *M*. The functor places  $_{f}(M)$  yielding a set is defined as follows:

(Def. 14) places  $_f(M)$  = the places of M.

The functor transitions  $_{f}(M)$  yielding a set is defined as follows:

(Def. 15) transitions  $_f(M)$  = the transitions of M.

The functor  $\operatorname{pre}_f(M)$  yielding a binary relation is defined by:

(Def. 16)  $\operatorname{pre}_f(M) = (\text{the flow relation of } M) \upharpoonright (\text{the transitions of } M).$ 

The functor  $post_f(M)$  yielding a binary relation is defined by:

(Def. 17)  $\text{post}_f(M) = (\text{the flow relation of } M) \subset [\text{the transitions of } M.$ 

Next we state three propositions:

- (27)  $\operatorname{pre}_f(M) \subseteq [:\operatorname{transitions}_f(M), \operatorname{places}_f(M):] \text{ and } \operatorname{post}_f(M) \subseteq [:\operatorname{transitions}_f(M), \operatorname{places}_f(M):].$
- (28) places  $_{f}(M)$  misses transitions  $_{f}(M)$ .
- (29)  $\operatorname{prox}_{f}(M) \subseteq [:\operatorname{Elements}(M), \operatorname{Elements}(M):] \text{ and } \operatorname{flow}_{f}(M) \subseteq [:\operatorname{Elements}(M), \operatorname{Elements}(M):].$

Let us consider *M*. The functor entrance f(M) yielding a binary relation is defined as follows:

(Def. 18) entrance<sub>*f*</sub>(*M*) = (the flow relation of *M*)  $\cong$  [the places of *M*  $\cup$  id<sub>the transitions of *M*.</sub>

The functor  $\operatorname{escape}_f(M)$  yielding a binary relation is defined as follows:

(Def. 19) escape  $_f(M) =$  (the flow relation of M) [(the places of M)  $\cup$  id<sub>the transitions of M.</sub>

Next we state four propositions:

- (30) escape  $_f(M) \subseteq [: \text{Elements}(M), \text{Elements}(M) :]$  and entrance  $_f(M) \subseteq [: \text{Elements}(M), \text{Elements}(M) :]$ .
- (31) domescape<sub>f</sub>(M)  $\subseteq$  Elements(M) and rng escape<sub>f</sub>(M)  $\subseteq$  Elements(M) and domentrance<sub>f</sub>(M)  $\subseteq$  Elements(M) and rng entrance<sub>f</sub>(M)  $\subseteq$  Elements(M).
- (32)  $\operatorname{escape}_f(M) \cdot \operatorname{escape}_f(M) = \operatorname{escape}_f(M)$  and  $\operatorname{escape}_f(M) \cdot \operatorname{entrance}_f(M) = \operatorname{escape}_f(M)$ and  $\operatorname{entrance}_f(M) \cdot \operatorname{entrance}_f(M) = \operatorname{entrance}_f(M)$  and  $\operatorname{entrance}_f(M) \cdot \operatorname{escape}_f(M) = \operatorname{entrance}_f(M)$ .
- (33)  $\operatorname{escape}_f(M) \cdot (\operatorname{escape}_f(M) \setminus \operatorname{id}_{\operatorname{Elements}(M)}) = \emptyset$  and  $\operatorname{entrance}_f(M) \cdot (\operatorname{entrance}_f(M) \setminus \operatorname{id}_{\operatorname{Elements}(M)}) = \emptyset$ .

Let us consider *M*. The functor  $adjac_f(M)$  yielding a binary relation is defined by the condition (Def. 20).

- (Def. 20)  $\operatorname{adjac}_f(M) = (\text{the flow relation of } M) \upharpoonright (\text{the transitions of } M) \cup (\text{the flow relation of } M)^{\sim} \upharpoonright \text{the transitions of } M \cup \operatorname{id}_{\operatorname{the transitions of } M}.$ 
  - We introduce circulation<sub>*f*</sub>(*M*) as a synonym of flow<sub>*f*</sub>(*M*). The following proposition is true
    - (34)  $\operatorname{adjac}_f(M) \cdot \operatorname{adjac}_f(M) = \operatorname{adjac}_f(M)$  and  $(\operatorname{adjac}_f(M) \setminus \operatorname{id}_{\operatorname{Elements}(M)}) \cdot \operatorname{adjac}_f(M) = \emptyset$  and  $\operatorname{adjac}_f(M) \cup (\operatorname{adjac}_f(M))^{\smile} \cup \operatorname{id}_{\operatorname{Elements}(M)} = \operatorname{circulation}_f(M) \cup (\operatorname{circulation}_f(M))^{\smile}$ .

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