## **Definitions of Petri Net. Part II**

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Summary. In the paper an equational definition of Petri net is given.

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The articles [3], [1], [4], and [2] provide the notation and terminology for this paper. In this paper *x*, *y*, *X*, *Y* are sets.

We introduce G-net structures which are extensions of 1-sorted structure and are systems  $\langle a \text{ carrier}, an \text{ entrance}, an \text{ escape } \rangle$ ,

where the carrier is a set and the entrance and the escape are binary relations.

Let *N* be a 1-sorted structure. The functor echaos(N) yields a set and is defined by:

(Def. 1) echaos(N) = (the carrier of N)  $\cup$  {the carrier of N}.

Let  $I_1$  be a G-net structure. We say that  $I_1$  is GG if and only if the conditions (Def. 2) are satisfied.

(Def. 2)(i) The entrance of  $I_1 \subseteq [:$  the carrier of  $I_1$ , the carrier of  $I_1$ :],

- (ii) the escape of  $I_1 \subseteq [:$  the carrier of  $I_1$ , the carrier of  $I_1:]$ ,
- (iii) (the entrance of  $I_1$ ) · (the entrance of  $I_1$ ) = the entrance of  $I_1$ ,
- (iv) (the entrance of  $I_1$ ) · (the escape of  $I_1$ ) = the entrance of  $I_1$ ,
- (v) (the escape of  $I_1$ ) · (the escape of  $I_1$ ) = the escape of  $I_1$ , and
- (vi) (the escape of  $I_1$ ) · (the entrance of  $I_1$ ) = the escape of  $I_1$ .

Let us observe that there exists a G-net structure which is GG. A G-net is a GG G-net structure.

Let  $I_1$  be a G-net structure. We say that  $I_1$  is EE if and only if the conditions (Def. 3) are satisfied.

(Def. 3)(i) (The entrance of  $I_1$ )  $\cdot$  ((the entrance of  $I_1$ )  $\setminus$  id<sub>the carrier of  $I_1$ ) =  $\emptyset$ , and</sub>

- (ii) (the escape of  $I_1$ )  $\cdot$  ((the escape of  $I_1$ )  $\setminus$  id<sub>the carrier of  $I_1$ ) =  $\emptyset$ .</sub>
- Let us mention that there exists a G-net structure which is EE.

Let us note that there exists a G-net structure which is strict, GG, and EE.

- An E-net is an EE GG G-net structure.
- In the sequel *N* denotes an E-net.

Next we state several propositions:

(1) Let *R*, *S* be binary relations. Then  $\langle X, R, S \rangle$  is an E-net if and only if the following conditions are satisfied:

 $R \subseteq [X, X]$  and  $S \subseteq [X, X]$  and  $R \cdot R = R$  and  $R \cdot S = R$  and  $S \cdot S = S$  and  $S \cdot R = S$  and  $R \cdot (R \setminus id_X) = \emptyset$  and  $S \cdot (S \setminus id_X) = \emptyset$ .

- (2)  $\langle X, \emptyset, \emptyset \rangle$  is an E-net.
- (3)  $\langle X, id_X, id_X \rangle$  is an E-net.
- (4)  $\langle \emptyset, \emptyset, \emptyset \rangle$  is an E-net.
- (8)<sup>1</sup>  $\langle X, \mathrm{id}_{X \setminus Y}, \mathrm{id}_{X \setminus Y} \rangle$  is an E-net.
- (9) echaos $(N) \neq \emptyset$ .

The strict E-net empty<sub>e</sub> is defined as follows:

(Def. 4) empty<sub>e</sub> =  $\langle \emptyset, \emptyset, \emptyset \rangle$ .

Let us consider X. The functor Tempty<sub>e</sub>(X) yielding a strict E-net is defined as follows:

(Def. 5) Tempty<sub>e</sub>(X) =  $\langle X, id_X, id_X \rangle$ .

The functor  $\text{Pempty}_{e}(X)$  yields a strict E-net and is defined as follows:

(Def. 6) Pempty<sub>e</sub>(X) =  $\langle X, \emptyset, \emptyset \rangle$ .

Next we state two propositions:

- (11)<sup>2</sup> The carrier of Tempty<sub>e</sub>(X) = X and the entrance of Tempty<sub>e</sub>(X) = id<sub>X</sub> and the escape of Tempty<sub>e</sub>(X) = id<sub>X</sub>.
- (12) The carrier of  $\text{Pempty}_e(X) = X$  and the entrance of  $\text{Pempty}_e(X) = \emptyset$  and the escape of  $\text{Pempty}_e(X) = \emptyset$ .

Let us consider x. The functor  $Psingle_e(x)$  yields a strict E-net and is defined by:

(Def. 7) Psingle<sub>e</sub>(x) =  $\langle \{x\}, id_{\{x\}}, id_{\{x\}} \rangle$ .

The functor  $Tsingle_{e}(x)$  yields a strict E-net and is defined as follows:

(Def. 8)  $\operatorname{Tsingle}_{e}(x) = \langle \{x\}, \emptyset, \emptyset \rangle.$ 

Next we state three propositions:

- (13) The carrier of  $Psingle_e(x) = \{x\}$  and the entrance of  $Psingle_e(x) = id_{\{x\}}$  and the escape of  $Psingle_e(x) = id_{\{x\}}$ .
- (14) The carrier of  $\text{Tsingle}_e(x) = \{x\}$  and the entrance of  $\text{Tsingle}_e(x) = \emptyset$  and the escape of  $\text{Tsingle}_e(x) = \emptyset$ .
- (15)  $\langle X \cup Y, id_X, id_X \rangle$  is an E-net.

Let us consider X, Y. The functor  $PTempty_e(X,Y)$  yielding a strict E-net is defined by:

(Def. 9) PTempty<sub>e</sub>(X, Y) =  $\langle X \cup Y, id_X, id_X \rangle$ .

One can prove the following propositions:

- (16)(i) (The entrance of *N*) \ id<sub>dom(the entrance of *N*) = (the entrance of *N*) \ id<sub>the carrier of *N*,</sub></sub>
- (ii) (the escape of N)  $\setminus$  id<sub>dom(the escape of N) = (the escape of N)  $\setminus$  id<sub>the carrier of N,</sub></sub>
- (iii) (the entrance of N)  $\setminus$  id<sub>rng(the entrance of N)</sub> = (the entrance of N)  $\setminus$  id<sub>the carrier of N</sub>, and
- (iv) (the escape of N) \ id<sub>rng(the escape of N) = (the escape of N) \ id<sub>the carrier of N</sub>.</sub>
- (17) CL(the entrance of N) = CL(the escape of N).

<sup>&</sup>lt;sup>1</sup> The propositions (5)–(7) have been removed.

<sup>&</sup>lt;sup>2</sup> The proposition (10) has been removed.

- (18)(i) rng (the entrance of N) = rng CL(the entrance of N),
- (ii) rng (the entrance of N) = dom CL(the entrance of N),
- (iii)  $\operatorname{rng}(\operatorname{the escape of} N) = \operatorname{rng} \operatorname{CL}(\operatorname{the escape of} N),$
- (iv)  $\operatorname{rng}(\operatorname{the escape of} N) = \operatorname{dom} \operatorname{CL}(\operatorname{the escape of} N), \text{ and }$
- (v)  $\operatorname{rng}(\text{the entrance of } N) = \operatorname{rng}(\text{the escape of } N).$
- (19)(i) dom(the entrance of N)  $\subseteq$  the carrier of N,
- (ii) rng (the entrance of N)  $\subseteq$  the carrier of N,
- (iii) dom(the escape of N)  $\subseteq$  the carrier of N, and
- (iv) rng (the escape of N)  $\subseteq$  the carrier of N.
- (20)(i) (The entrance of *N*)  $\cdot$  ((the escape of *N*)  $\setminus$  id<sub>the carrier of *N*</sub>) =  $\emptyset$ , and
- (ii) (the escape of N)  $\cdot$  ((the entrance of N)  $\setminus$  id<sub>the carrier of N</sub>) =  $\emptyset$ .
- (21)(i) ((The entrance of N) \ id<sub>the carrier of N</sub>) · ((the entrance of N) \ id<sub>the carrier of N</sub>) =  $\emptyset$ ,
- (ii) ((the escape of N) \ id\_{the carrier of N})  $\cdot$  ((the escape of N) \ id\_{the carrier of N}) =  $\emptyset$ ,
- (iii) ((the entrance of N) \ id<sub>the carrier of N</sub>)  $\cdot$  ((the escape of N) \ id<sub>the carrier of N</sub>) =  $\emptyset$ , and
- (iv) ((the escape of *N*) \ id<sub>the carrier of *N*</sub>)  $\cdot$  ((the entrance of *N*) \ id<sub>the carrier of *N*</sub>) =  $\emptyset$ .

Let us consider N. The functor  $Places_e(N)$  yielding a set is defined by:

(Def. 10) Places<sub>e</sub>(
$$N$$
) = rng(the entrance of  $N$ ).

Let us consider N. The functor  $Transitions_e(N)$  yields a set and is defined by:

(Def. 11) Transitions<sub>e</sub>(
$$N$$
) = (the carrier of  $N$ ) \ Places<sub>e</sub>( $N$ ).

Next we state three propositions:

- (22) Places<sub>e</sub>(N) misses Transitions<sub>e</sub>(N).
- (23) If  $\langle x, y \rangle \in$  the entrance of *N* or  $\langle x, y \rangle \in$  the escape of *N* and if  $x \neq y$ , then  $x \in$  Transitions<sub>e</sub>(*N*) and  $y \in$  Places<sub>e</sub>(*N*).
- (24) (The entrance of N) \ id<sub>the carrier of  $N \subseteq [:$  Transitions<sub>e</sub>(N), Places<sub>e</sub>(N):] and (the escape of N) \ id<sub>the carrier of  $N \subseteq [:$  Transitions<sub>e</sub>(N), Places<sub>e</sub>(N):].</sub></sub>

Let us consider N. The functor  $Flow_e(N)$  yielding a binary relation is defined by:

(Def. 12) Flow<sub>e</sub>(N) = ((the entrance of N)  $\subset \cup$  the escape of N)  $\setminus$  id<sub>the carrier of N</sub>.

Next we state the proposition

(25)  $\operatorname{Flow}_{e}(N) \subseteq [\operatorname{Places}_{e}(N), \operatorname{Transitions}_{e}(N)] \cup [\operatorname{Transitions}_{e}(N), \operatorname{Places}_{e}(N)].$ 

Let us consider N. We introduce  $places_e(N)$  as a synonym of  $Places_e(N)$ . We introduce transitions<sub>e</sub>(N) as a synonym of Transitions<sub>e</sub>(N).

Let us consider N. The functor  $\operatorname{pre}_{e}(N)$  yields a binary relation and is defined by:

(Def. 15)<sup>3</sup>  $\operatorname{pre}_{e}(N) = (\text{the entrance of } N) \setminus \operatorname{id}_{\operatorname{the carrier of } N}.$ 

The functor  $post_e(N)$  yielding a binary relation is defined as follows:

(Def. 16)  $\text{post}_e(N) = (\text{the escape of } N) \setminus \text{id}_{\text{the carrier of } N}.$ 

We now state the proposition

<sup>&</sup>lt;sup>3</sup> The definitions (Def. 13) and (Def. 14) have been removed.

 $(28)^4$  pre<sub>e</sub>(N)  $\subseteq$  [: transitions<sub>e</sub>(N), places<sub>e</sub>(N) :] and post<sub>e</sub>(N)  $\subseteq$  [: transitions<sub>e</sub>(N), places<sub>e</sub>(N) :].

Let us consider N. The functor shore<sub>e</sub>(N) yields a set and is defined by:

(Def. 17) shore<sub>e</sub>(N) = the carrier of N.

The functor  $\operatorname{prox}_{\rho}(N)$  yields a binary relation and is defined by:

(Def. 18)  $\operatorname{prox}_{e}(N) = ((\text{the entrance of } N) \cup (\text{the escape of } N))^{\smile}.$ 

The functor  $flow_e(N)$  yields a binary relation and is defined by:

(Def. 19) flow<sub>e</sub>(N) = (the entrance of N)  $\smile \cup$  the escape of  $N \cup id_{\text{the carrier of } N}$ .

Next we state several propositions:

- (29)  $\operatorname{prox}_{e}(N) \subseteq [:\operatorname{shore}_{e}(N), \operatorname{shore}_{e}(N):] \text{ and } \operatorname{flow}_{e}(N) \subseteq [:\operatorname{shore}_{e}(N), \operatorname{shore}_{e}(N):].$
- (30)  $\operatorname{prox}_{e}(N) \cdot \operatorname{prox}_{e}(N) = \operatorname{prox}_{e}(N)$  and  $(\operatorname{prox}_{e}(N) \setminus \operatorname{id}_{\operatorname{shore}_{e}(N)}) \cdot \operatorname{prox}_{e}(N) = \emptyset$  and  $\operatorname{prox}_{e}(N) \cup (\operatorname{prox}_{e}(N))^{\smile} \cup \operatorname{id}_{\operatorname{shore}_{e}(N)} = \operatorname{flow}_{e}(N) \cup (\operatorname{flow}_{e}(N))^{\smile}.$
- (31)(i)  $\operatorname{id}_{(\text{the carrier of }N)\setminus\operatorname{rng}(\text{the escape of }N)} \cdot ((\text{the escape of }N) \setminus \operatorname{id}_{\text{the carrier of }N}) = (\text{the escape of }N) \setminus \operatorname{id}_{\text{the carrier of }N}, \text{ and}$
- (ii)  $\operatorname{id}_{(\text{the carrier of }N)\setminus \operatorname{rng}(\text{the entrance of }N)} \cdot ((\text{the entrance of }N) \setminus \operatorname{id}_{\operatorname{the carrier of }N}) = (\text{the entrance of }N) \setminus \operatorname{id}_{\operatorname{the carrier of }N}.$
- (32)(i) ((The escape of *N*) \ id<sub>the carrier of *N*</sub>)  $\cdot$  ((the escape of *N*) \ id<sub>the carrier of *N*</sub>) =  $\emptyset$ ,
- (ii) ((the entrance of *N*) \ id<sub>the carrier of *N*</sub>)  $\cdot$  ((the entrance of *N*) \ id<sub>the carrier of *N*</sub>) =  $\emptyset$ ,
- (iii) ((the escape of *N*) \ id<sub>the carrier of *N*</sub>)  $\cdot$  ((the entrance of *N*) \ id<sub>the carrier of *N*</sub>) =  $\emptyset$ , and
- (iv) ((the entrance of *N*) \  $id_{the \ carrier \ of \ N}$ )  $\cdot$  ((the escape of *N*) \  $id_{the \ carrier \ of \ N}$ ) =  $\emptyset$ .
- (33)(i) ((The escape of N) \ id\_{the carrier of N})  $\sim \cdot$  ((the escape of N) \ id\_{the carrier of N})  $\sim = \emptyset$ , and
- (ii) ((the entrance of N) \ id<sub>the carrier of N</sub>)  $\simeq$  · ((the entrance of N) \ id<sub>the carrier of N</sub>)  $\simeq \emptyset$ .
- (34)(i) ((The escape of N) \ id\_{the carrier of N})  $\sim \cdot (id_{(the carrier of N) \setminus rng(the escape of N)}) = ((the escape of N) \setminus id_{the carrier of N})$ , and
- (ii) ((the entrance of N) \  $id_{the \ carrier \ of \ N}$ )  $\sim \cdot (id_{(the \ carrier \ of \ N) \setminus rng(the \ entrance \ of \ N)})^{\sim} = ((the \ entrance \ of \ N) \setminus id_{the \ carrier \ of \ N})^{\sim}.$
- (35)(i) ((The escape of N) \ id\_{the carrier of N})  $\cdot$  id\_{(the carrier of N) \rng(the escape of N)} =  $\emptyset$ , and
- (ii) ((the entrance of N) \ id\_{the carrier of N}) \cdot id\_{(the carrier of N) \setminus rng(the entrance of N)} =  $\emptyset$ .
- (36)(i)  $\operatorname{id}_{\operatorname{(the carrier of N)} \operatorname{rng}(\operatorname{the escape of N})} \cdot ((\operatorname{the escape of N}) \setminus \operatorname{id}_{\operatorname{the carrier of N}})^{\smile} = \emptyset$ , and
- (ii)  $\operatorname{id}_{(\text{the carrier of }N)\setminus\operatorname{rng}(\text{the entrance of }N)} \cdot ((\text{the entrance of }N) \setminus \operatorname{id}_{\text{the carrier of }N})^{\smile} = \emptyset.$

Let us consider N. We introduce  $\operatorname{support}_e(N)$  as a synonym of  $\operatorname{shore}_e(N)$ . Let us consider N. The functor  $\operatorname{entrance}_e(N)$  yields a binary relation and is defined by:

 $(\text{Def. 21})^5 \quad \text{entrance}_e(N) = ((\text{the escape of } N) \setminus \text{id}_{\text{the carrier of } N})^{\smile} \cup \text{id}_{(\text{the carrier of } N) \setminus \text{rng}(\text{the escape of } N)}.$ 

The functor  $escape_e(N)$  yields a binary relation and is defined by:

- (Def. 22)  $\operatorname{escape}_{e}(N) = ((\operatorname{the entrance of } N) \setminus \operatorname{id}_{\operatorname{the carrier of } N})^{\smile} \cup \operatorname{id}_{\operatorname{(the carrier of } N) \setminus \operatorname{rng}(\operatorname{the entrance of } N)}$ . We now state two propositions:
  - (37) entrance<sub>e</sub>(N) entrance<sub>e</sub>(N) = entrance<sub>e</sub>(N) and entrance<sub>e</sub>(N) escape<sub>e</sub>(N) = entrance<sub>e</sub>(N) and escape<sub>e</sub>(N) escape<sub>e</sub>(N) = escape<sub>e</sub>(N).

<sup>&</sup>lt;sup>4</sup> The propositions (26) and (27) have been removed.

<sup>&</sup>lt;sup>5</sup> The definition (Def. 20) has been removed.

(38) entrance<sub>e</sub>(N) · (entrance<sub>e</sub>(N) \ id<sub>support<sub>e</sub>(N)</sub>) =  $\emptyset$  and escape<sub>e</sub>(N) · (escape<sub>e</sub>(N) \ id<sub>support<sub>e</sub>(N)</sub>) =  $\emptyset$ .

Let us consider N. We introduce stanchion<sub>e</sub>(N) as a synonym of shore<sub>e</sub>(N). Let us consider N. The functor  $adjac_e(N)$  yields a binary relation and is defined as follows:

- (Def. 24)<sup>6</sup>  $\operatorname{adjac}_{e}(N) = (((\text{the entrance of } N) \cup (\text{the escape of } N)) \setminus \operatorname{id}_{\operatorname{the carrier of } N}) \cup \operatorname{id}_{\operatorname{(the carrier of } N) \setminus \operatorname{rng}(\operatorname{the entrance of } N)})$ 
  - We introduce circulation<sub>e</sub>(N) as a synonym of  $flow_e(N)$ . One can prove the following two propositions:
    - (39)  $\operatorname{adjac}_{e}(N) \subseteq [\operatorname{stanchion}_{e}(N), \operatorname{stanchion}_{e}(N):]$  and  $\operatorname{circulation}_{e}(N) \subseteq [\operatorname{stanchion}_{e}(N), \operatorname{stanchion}_{e}(N):]$ .
    - (40)  $\operatorname{adjac}_{e}(N) \cdot \operatorname{adjac}_{e}(N) = \operatorname{adjac}_{e}(N)$  and  $(\operatorname{adjac}_{e}(N) \setminus \operatorname{id}_{\operatorname{stanchion}_{e}(N)}) \cdot \operatorname{adjac}_{e}(N) = \emptyset$  and  $\operatorname{adjac}_{e}(N) \cup (\operatorname{adjac}_{e}(N))^{\smile} \cup \operatorname{id}_{\operatorname{stanchion}_{e}(N)} = \operatorname{circulation}_{e}(N) \cup (\operatorname{circulation}_{e}(N))^{\smile}$ .

Let N be an E-net. We introduce transitions<sub>s</sub>(N) as a synonym of  $Places_e(N)$ . We introduce  $places_s(N)$  as a synonym of  $Transitions_e(N)$ . We introduce  $carrier_s(N)$  as a synonym of shore<sub>e</sub>(N). We introduce  $enter_s(N)$  as a synonym of  $entrance_e(N)$ . We introduce  $exit_s(N)$  as a synonym of  $escape_e(N)$ . We introduce  $prox_s(N)$  as a synonym of  $adjac_e(N)$ .

In the sequel *N* is an E-net.

One can prove the following proposition

(41) ((The entrance of N)  $\setminus$  id<sub>the carrier of N)  $\cong$  [:Places<sub>e</sub>(N), Transitions<sub>e</sub>(N):] and ((the escape of N)  $\setminus$  id<sub>the carrier of N)  $\cong$  [:Places<sub>e</sub>(N), Transitions<sub>e</sub>(N):].</sub></sub>

Let N be a G-net structure. The functor  $\operatorname{pre}_{\mathfrak{s}}(N)$  yielding a binary relation is defined as follows:

(Def. 25)  $\operatorname{pre}_{s}(N) = ((\operatorname{the escape of } N) \setminus \operatorname{id}_{\operatorname{the carrier of } N})^{\smile}.$ 

The functor  $post_s(N)$  yielding a binary relation is defined as follows:

(Def. 26)  $\operatorname{post}_{s}(N) = ((\text{the entrance of } N) \setminus \operatorname{id}_{\operatorname{the carrier of } N})^{\smile}.$ 

The following proposition is true

(42)  $\text{post}_{s}(N) \subseteq [:\text{transitions}_{s}(N), \text{places}_{s}(N):] \text{ and } \text{pre}_{s}(N) \subseteq [:\text{transitions}_{s}(N), \text{places}_{s}(N):].$ 

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<sup>&</sup>lt;sup>6</sup> The definition (Def. 23) has been removed.