

Logical Equivalence of Formulae¹

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The articles [8], [9], [7], [1], [3], [2], [6], [4], and [5] provide the notation and terminology for this paper.

For simplicity, we adopt the following rules: p, q, r, s, p_1, q_1 denote elements of CQC-WFF, X, Y, Z, X_1, X_2 denote subsets of CQC-WFF, h denotes a formula, and x, y denote bound variables.

The following four propositions are true:

- (1) If $p \in X$, then $X \vdash p$.
- (2) If $X \subseteq \text{Cn } Y$, then $\text{Cn } X \subseteq \text{Cn } Y$.
- (3) If $X \vdash p$ and $\{p\} \vdash q$, then $X \vdash q$.
- (4) If $X \vdash p$ and $X \subseteq Y$, then $Y \vdash p$.

Let p, q be elements of CQC-WFF. The predicate $p \vdash q$ is defined by:

(Def. 1) $\{p\} \vdash q$.

The following propositions are true:

- (5) $p \vdash p$.
- (6) If $p \vdash q$ and $q \vdash r$, then $p \vdash r$.

Let X, Y be subsets of CQC-WFF. The predicate $X \vdash Y$ is defined by:

(Def. 2) For every element p of CQC-WFF such that $p \in Y$ holds $X \vdash p$.

Next we state several propositions:

- (7) $X \vdash Y$ iff $Y \subseteq \text{Cn } X$.
- (8) $X \vdash X$.
- (9) If $X \vdash Y$ and $Y \vdash Z$, then $X \vdash Z$.
- (10) $X \vdash \{p\}$ iff $X \vdash p$.
- (11) $\{p\} \vdash \{q\}$ iff $p \vdash q$.

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(12) If $X \subseteq Y$, then $Y \vdash X$.

(13) $X \vdash \text{Taut}$.

(14) $\emptyset_{\text{CQC-WFF}} \vdash \text{Taut}$.

Let X be a subset of CQC-WFF. The predicate $\vdash X$ is defined as follows:

(Def. 3) For every element p of CQC-WFF such that $p \in X$ holds $\vdash p$.

The following propositions are true:

(15) $\vdash X$ iff $\emptyset_{\text{CQC-WFF}} \vdash X$.

(16) $\vdash \text{Taut}$.

(17) $\vdash X$ iff $X \subseteq \text{Taut}$.

Let us consider X, Y . The predicate $X \vdash Y$ is defined by:

(Def. 4) For every p holds $X \vdash p$ iff $Y \vdash p$.

Let us notice that the predicate $X \vdash Y$ is reflexive and symmetric.

The following propositions are true:

(18) $X \vdash Y$ iff $X \vdash Y$ and $Y \vdash X$.

(19) If $X \vdash Y$ and $Y \vdash Z$, then $X \vdash Z$.

(20) $X \vdash Y$ iff $\text{Cn}X = \text{Cn}Y$.

(21) $\text{Cn}X \cup \text{Cn}Y \subseteq \text{Cn}(X \cup Y)$.

(22) $\text{Cn}(X \cup Y) = \text{Cn}(\text{Cn}X \cup \text{Cn}Y)$.

(23) $X \vdash \text{Cn}X$.

(24) $X \cup Y \vdash \text{Cn}X \cup \text{Cn}Y$.

(25) If $X_1 \vdash X_2$, then $X_1 \cup Y \vdash X_2 \cup Y$.

(26) If $X_1 \vdash X_2$ and $X_1 \cup Y \vdash Z$, then $X_2 \cup Y \vdash Z$.

(27) If $X_1 \vdash X_2$ and $Y \vdash X_1$, then $Y \vdash X_2$.

Let p, q be elements of CQC-WFF. The predicate $p \vdash q$ is defined by:

(Def. 5) $p \vdash q$ and $q \vdash p$.

Let us notice that the predicate $p \vdash q$ is reflexive and symmetric.

The following propositions are true:

(28) If $p \vdash q$ and $q \vdash r$, then $p \vdash r$.

(29) $p \vdash q$ iff $\{p\} \vdash \{q\}$.

(30) If $p \vdash q$ and $X \vdash p$, then $X \vdash q$.

(31) $\{p, q\} \vdash \{p \wedge q\}$.

(32) $p \wedge q \vdash q \wedge p$.

(33) $X \vdash p \wedge q$ iff $X \vdash p$ and $X \vdash q$.

(34) If $p \vdash q$ and $r \vdash s$, then $p \wedge r \vdash q \wedge s$.

(35) $X \vdash \forall_x p$ iff $X \vdash p$.

$$(36) \quad \forall_x p \vdash\vdash p.$$

$$(37) \quad \text{If } p \vdash\vdash q, \text{ then } \forall_x p \vdash\vdash \forall_y q.$$

Let p, q be elements of CQC-WFF. We say that p is an universal closure of q if and only if the conditions (Def. 6) are satisfied.

(Def. 6)(i) p is closed, and

- (ii) there exists a natural number n such that $1 \leq n$ and there exists a finite sequence L such that $\text{len } L = n$ and $L(1) = q$ and $L(n) = p$ and for every natural number k such that $1 \leq k$ and $k < n$ there exists a bound variable x and there exists an element r of CQC-WFF such that $r = L(k)$ and $L(k+1) = \forall_x r$.

One can prove the following propositions:

$$(38) \quad \text{If } p \text{ is an universal closure of } q, \text{ then } p \vdash\vdash q.$$

$$(39) \quad \text{If } \vdash p \Rightarrow q, \text{ then } p \vdash q.$$

$$(40) \quad \text{If } X \vdash p \Rightarrow q, \text{ then } X \cup \{p\} \vdash q.$$

$$(41) \quad \text{If } p \text{ is closed and } p \vdash q, \text{ then } \vdash p \Rightarrow q.$$

$$(42) \quad \text{If } p_1 \text{ is an universal closure of } p, \text{ then } X \cup \{p\} \vdash q \text{ iff } X \vdash p_1 \Rightarrow q.$$

$$(43) \quad \text{If } p \text{ is closed and } p \vdash q, \text{ then } \neg q \vdash \neg p.$$

$$(44) \quad \text{If } p \text{ is closed and } X \cup \{p\} \vdash q, \text{ then } X \cup \{\neg q\} \vdash \neg p.$$

$$(45) \quad \text{If } p \text{ is closed and } \neg p \vdash \neg q, \text{ then } q \vdash p.$$

$$(46) \quad \text{If } p \text{ is closed and } X \cup \{\neg p\} \vdash \neg q, \text{ then } X \cup \{q\} \vdash p.$$

$$(47) \quad \text{If } p \text{ is closed and } q \text{ is closed, then } p \vdash q \text{ iff } \neg q \vdash \neg p.$$

$$(48) \quad \text{If } p_1 \text{ is an universal closure of } p \text{ and } q_1 \text{ is an universal closure of } q, \text{ then } p \vdash q \text{ iff } \neg q_1 \vdash \neg p_1.$$

$$(49) \quad \text{If } p_1 \text{ is an universal closure of } p \text{ and } q_1 \text{ is an universal closure of } q, \text{ then } p \vdash\vdash q \text{ iff } \neg p_1 \vdash\vdash \neg q_1.$$

Let p, q be elements of CQC-WFF. The predicate $p \equiv q$ is defined as follows:

(Def. 7) $\vdash p \Leftrightarrow q$.

Let us notice that the predicate $p \equiv q$ is reflexive and symmetric.

One can prove the following propositions:

$$(50) \quad p \equiv q \text{ iff } \vdash p \Rightarrow q \text{ and } \vdash q \Rightarrow p.$$

$$(51) \quad \text{If } p \equiv q \text{ and } q \equiv r, \text{ then } p \equiv r.$$

$$(52) \quad \text{If } p \equiv q, \text{ then } p \vdash\vdash q.$$

$$(53) \quad p \equiv q \text{ iff } \neg p \equiv \neg q.$$

$$(54) \quad \text{If } p \equiv q \text{ and } r \equiv s, \text{ then } p \wedge r \equiv q \wedge s.$$

$$(55) \quad \text{If } p \equiv q \text{ and } r \equiv s, \text{ then } p \Rightarrow r \equiv q \Rightarrow s.$$

$$(56) \quad \text{If } p \equiv q \text{ and } r \equiv s, \text{ then } p \vee r \equiv q \vee s.$$

$$(57) \quad \text{If } p \equiv q \text{ and } r \equiv s, \text{ then } p \Leftrightarrow r \equiv q \Leftrightarrow s.$$

$$(58) \quad \text{If } p \equiv q, \text{ then } \forall_x p \equiv \forall_x q.$$

- (59) If $p \equiv q$, then $\exists_x p \equiv \exists_x q$.
- (61)¹ Let k be a natural number, l be a list of variables of the length k , a be a free variable, and x be a bound variable. Then $\text{snb}(l) \subseteq \text{snb}(l[a \mapsto x])$.
- (62) Let k be a natural number, l be a list of variables of the length k , a be a free variable, and x be a bound variable. Then $\text{snb}(l[a \mapsto x]) \subseteq \text{snb}(l) \cup \{x\}$.
- (63) For every h holds $\text{snb}(h) \subseteq \text{snb}(h(x))$.
- (64) For every h holds $\text{snb}(h(x)) \subseteq \text{snb}(h) \cup \{x\}$.
- (65) If $p = h(x)$ and $x \neq y$ and $y \notin \text{snb}(h)$, then $y \notin \text{snb}(p)$.
- (66) If $p = h(x)$ and $q = h(y)$ and $x \notin \text{snb}(h)$ and $y \notin \text{snb}(h)$, then $\forall_x p \equiv \forall_y q$.

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¹ The proposition (60) has been removed.