

Propositional Calculus for Boolean Valued Functions.

Part VI

Shunichi Kobayashi
Shinshu University
Nagano

Summary. In this paper, we proved some elementary propositional calculus formulae for Boolean valued functions.

MML Identifier: BVFUNC10.

WWW: <http://mizar.org/JFM/Vol11/bvfunc10.html>

The articles [3], [5], [4], [2], and [1] provide the notation and terminology for this paper.

In this paper Y denotes a non empty set.

One can prove the following propositions:

- (1) For all elements a, b, c of Boolean^Y holds $a \wedge b \vee b \wedge c \vee c \wedge a = (a \vee b) \wedge (b \vee c) \wedge (c \vee a)$.
- (2) For all elements a, b, c of Boolean^Y holds $a \wedge \neg b \vee b \wedge \neg c \vee c \wedge \neg a = b \wedge \neg a \vee c \wedge \neg b \vee a \wedge \neg c$.
- (3) For all elements a, b, c of Boolean^Y holds $(a \vee \neg b) \wedge (b \vee \neg c) \wedge (c \vee \neg a) = (b \vee \neg a) \wedge (c \vee \neg b) \wedge (a \vee \neg c)$.
- (4) For all elements a, b, c of Boolean^Y such that $c \Rightarrow a = \text{true}(Y)$ and $c \Rightarrow b = \text{true}(Y)$ holds $c \Rightarrow a \vee b = \text{true}(Y)$.
- (5) For all elements a, b, c of Boolean^Y such that $a \Rightarrow c = \text{true}(Y)$ and $b \Rightarrow c = \text{true}(Y)$ holds $a \wedge b \Rightarrow c = \text{true}(Y)$.
- (6) For all elements $a_1, a_2, b_1, b_2, c_1, c_2$ of Boolean^Y holds $(a_1 \Rightarrow a_2) \wedge (b_1 \Rightarrow b_2) \wedge (c_1 \Rightarrow c_2) \wedge (a_1 \vee b_1 \vee c_1) \in a_2 \vee b_2 \vee c_2$.
- (7) For all elements a_1, a_2, b_1, b_2 of Boolean^Y holds $(a_1 \Rightarrow b_1) \wedge (a_2 \Rightarrow b_2) \wedge (a_1 \vee a_2) \wedge \neg(b_1 \wedge b_2) = (b_1 \Rightarrow a_1) \wedge (b_2 \Rightarrow a_2) \wedge (b_1 \vee b_2) \wedge \neg(a_1 \wedge a_2)$.
- (8) For all elements a, b, c, d of Boolean^Y holds $(a \wedge b) \wedge (c \vee d) = a \wedge c \vee a \wedge d \vee b \wedge c \vee b \wedge d$.
- (9) For all elements a_1, a_2, b_1, b_2, b_3 of Boolean^Y holds $a_1 \wedge a_2 \vee b_1 \wedge b_2 \wedge b_3 = (a_1 \vee b_1) \wedge (a_1 \vee b_2) \wedge (a_1 \vee b_3) \wedge (a_2 \vee b_1) \wedge (a_2 \vee b_2) \wedge (a_2 \vee b_3)$.
- (10) For all elements a, b, c, d of Boolean^Y holds $(a \Rightarrow b) \wedge (b \Rightarrow c) \wedge (c \Rightarrow d) = (a \Rightarrow b \wedge c \wedge d) \wedge (b \Rightarrow c \wedge d) \wedge (c \Rightarrow d)$.
- (11) For all elements a, b, c, d of Boolean^Y holds $(a \Rightarrow c) \wedge (b \Rightarrow d) \wedge (a \vee b) \in c \vee d$.
- (12) For all elements a, b, c of Boolean^Y holds $(a \wedge b \Rightarrow \neg c) \wedge a \wedge c \in \neg b$.

- (13) For all elements $a_1, a_2, a_3, b_1, b_2, b_3$ of Boolean^Y holds $a_1 \wedge a_2 \wedge a_3 \Rightarrow b_1 \vee b_2 \vee b_3 = \neg b_1 \wedge \neg b_2 \wedge a_3 \Rightarrow \neg a_1 \vee \neg a_2 \vee b_3$.
- (14) For all elements a, b, c of Boolean^Y holds $(a \Rightarrow b) \wedge (b \Rightarrow c) \wedge (c \Rightarrow a) = a \wedge b \wedge c \vee \neg a \wedge \neg b \wedge \neg c$.
- (15) For all elements a, b, c of Boolean^Y holds $(a \Rightarrow b) \wedge (b \Rightarrow c) \wedge (c \Rightarrow a) \wedge (a \vee b \vee c) = a \wedge b \wedge c$.
- (16) For all elements a, b, c of Boolean^Y holds $(a \vee b) \wedge (b \vee c) \wedge (c \vee a) \wedge \neg(a \wedge b \wedge c) = \neg a \wedge b \wedge c \vee a \wedge \neg b \wedge c \vee a \wedge b \wedge \neg c$.
- (17) For all elements a, b, c of Boolean^Y holds $(a \Rightarrow b) \wedge (b \Rightarrow c) \in a \Rightarrow b \wedge c$.
- (18) For all elements a, b, c of Boolean^Y holds $(a \Rightarrow b) \wedge (b \Rightarrow c) \in a \vee b \Rightarrow c$.
- (19) For all elements a, b, c of Boolean^Y holds $(a \Rightarrow b) \wedge (b \Rightarrow c) \in a \Rightarrow b \vee c$.
- (20) For all elements a, b, c of Boolean^Y holds $(a \Rightarrow b) \wedge (b \Rightarrow c) \in a \Rightarrow b \vee \neg c$.
- (21) For all elements a, b, c of Boolean^Y holds $(a \Rightarrow b) \wedge (b \Rightarrow c) \in b \Rightarrow c \vee a$.
- (22) For all elements a, b, c of Boolean^Y holds $(a \Rightarrow b) \wedge (b \Rightarrow c) \in b \Rightarrow c \vee \neg a$.
- (23) For all elements a, b, c of Boolean^Y holds $(a \Rightarrow b) \wedge (b \Rightarrow c) \in (a \Rightarrow b) \wedge (b \Rightarrow c \vee a)$.
- (24) For all elements a, b, c of Boolean^Y holds $(a \Rightarrow b) \wedge (b \Rightarrow c) \in (a \Rightarrow b \vee \neg c) \wedge (b \Rightarrow c)$.
- (25) For all elements a, b, c of Boolean^Y holds $(a \Rightarrow b) \wedge (b \Rightarrow c) \in (a \Rightarrow b \vee c) \wedge (b \Rightarrow c \vee a)$.
- (26) For all elements a, b, c of Boolean^Y holds $(a \Rightarrow b) \wedge (b \Rightarrow c) \in (a \Rightarrow b \vee \neg c) \wedge (b \Rightarrow c \vee a)$.
- (27) For all elements a, b, c of Boolean^Y holds $(a \Rightarrow b) \wedge (b \Rightarrow c) \in (a \Rightarrow b \vee \neg c) \wedge (b \Rightarrow c \vee \neg a)$.

REFERENCES

- [1] Shunichi Kobayashi and Kui Jia. A theory of Boolean valued functions and partitions. *Journal of Formalized Mathematics*, 10, 1998. http://mizar.org/JFM/Vol10/bvfunc_1.html.
- [2] Andrzej Trybulec. Function domains and Fränkel operator. *Journal of Formalized Mathematics*, 2, 1990. <http://mizar.org/JFM/Vol2/fraenkel.html>.
- [3] Zinaida Trybulec. Properties of subsets. *Journal of Formalized Mathematics*, 1, 1989. http://mizar.org/JFM/Vol1/subset_1.html.
- [4] Edmund Woronowicz. Interpretation and satisfiability in the first order logic. *Journal of Formalized Mathematics*, 2, 1990. http://mizar.org/JFM/Vol2/valuat_1.html.
- [5] Edmund Woronowicz. Many-argument relations. *Journal of Formalized Mathematics*, 2, 1990. <http://mizar.org/JFM/Vol2/margrell.html>.

Received July 14, 1999

Published January 2, 2004
