

On the Instructions of SCM¹

Artur Korniłowicz
University of Białystok

MML Identifier: AMI_6.

WWW: http://mizar.org/JFM/Vol13/ami_6.html

The articles [17], [23], [18], [14], [24], [7], [8], [2], [3], [19], [22], [1], [9], [5], [10], [15], [4], [6], [12], [13], [20], [16], [21], and [11] provide the notation and terminology for this paper.

For simplicity, we follow the rules: a, b denote data-locations, i_1, i_2, i_3 denote instruction-locations of **SCM**, s_1, s_2 denote states of **SCM**, T denotes an instruction type of **SCM**, and k denotes a natural number.

The following propositions are true:

- (1) $a \notin$ the instruction locations of **SCM**.
- (2) $\text{Data-Loc}_{\text{SCM}} \neq$ the instruction locations of **SCM**.
- (3) For every object o of **SCM** holds $o = \text{IC}_{\text{SCM}}$ or $o \in$ the instruction locations of **SCM** or o is a data-location.
- (4) If $i_2 \neq i_3$, then $\text{Next}(i_2) \neq \text{Next}(i_3)$.
- (5) If s_1 and s_2 are equal outside the instruction locations of **SCM**, then $s_1(a) = s_2(a)$.
- (6) Let N be a set with non empty elements, S be a realistic IC-Ins-separated definite non empty non void AMI over N , t, u be states of S , i_1 be an instruction-location of S , e be an element of $\text{ObjectKind}(\text{IC}_S)$, and I be an element of $\text{ObjectKind}(i_1)$. If $e = i_1$ and $u = t + \cdot [\text{IC}_S \mapsto e, i_1 \mapsto I]$, then $u(i_1) = I$ and $\text{IC}_u = i_1$ and $\text{IC}_{\text{Following}(u)} = (\text{Exec}(u(\text{IC}_u), u))(\text{IC}_S)$.
- (7) $\text{AddressPart}(\text{halt}_{\text{SCM}}) = \emptyset$.
- (8) $\text{AddressPart}(a:=b) = \langle a, b \rangle$.
- (9) $\text{AddressPart}(\text{AddTo}(a, b)) = \langle a, b \rangle$.
- (10) $\text{AddressPart}(\text{SubFrom}(a, b)) = \langle a, b \rangle$.
- (11) $\text{AddressPart}(\text{MultBy}(a, b)) = \langle a, b \rangle$.
- (12) $\text{AddressPart}(\text{Divide}(a, b)) = \langle a, b \rangle$.
- (13) $\text{AddressPart}(\text{goto } i_2) = \langle i_2 \rangle$.
- (14) $\text{AddressPart}(\text{if } a = 0 \text{ goto } i_2) = \langle i_2, a \rangle$.
- (15) $\text{AddressPart}(\text{if } a > 0 \text{ goto } i_2) = \langle i_2, a \rangle$.

¹This work has been partially supported by TYPES grant IST-1999-29001.

(16) If $T = 0$, then $\text{AddressParts } T = \{0\}$.

Let us consider T . Observe that $\text{AddressParts } T$ is non empty.

We now state a number of propositions:

(17) If $T = 1$, then $\text{dom } \prod_{\text{AddressParts } T} = \{1, 2\}$.

(18) If $T = 2$, then $\text{dom } \prod_{\text{AddressParts } T} = \{1, 2\}$.

(19) If $T = 3$, then $\text{dom } \prod_{\text{AddressParts } T} = \{1, 2\}$.

(20) If $T = 4$, then $\text{dom } \prod_{\text{AddressParts } T} = \{1, 2\}$.

(21) If $T = 5$, then $\text{dom } \prod_{\text{AddressParts } T} = \{1, 2\}$.

(22) If $T = 6$, then $\text{dom } \prod_{\text{AddressParts } T} = \{1\}$.

(23) If $T = 7$, then $\text{dom } \prod_{\text{AddressParts } T} = \{1, 2\}$.

(24) If $T = 8$, then $\text{dom } \prod_{\text{AddressParts } T} = \{1, 2\}$.

(25) $\prod_{\text{AddressParts } \text{InsCode}(a:=b)}(1) = \text{Data-Loc}_{\text{SCM}}$.

(26) $\prod_{\text{AddressParts } \text{InsCode}(a:=b)}(2) = \text{Data-Loc}_{\text{SCM}}$.

(27) $\prod_{\text{AddressParts } \text{InsCode}(\text{AddTo}(a,b))}(1) = \text{Data-Loc}_{\text{SCM}}$.

(28) $\prod_{\text{AddressParts } \text{InsCode}(\text{AddTo}(a,b))}(2) = \text{Data-Loc}_{\text{SCM}}$.

(29) $\prod_{\text{AddressParts } \text{InsCode}(\text{SubFrom}(a,b))}(1) = \text{Data-Loc}_{\text{SCM}}$.

(30) $\prod_{\text{AddressParts } \text{InsCode}(\text{SubFrom}(a,b))}(2) = \text{Data-Loc}_{\text{SCM}}$.

(31) $\prod_{\text{AddressParts } \text{InsCode}(\text{MultBy}(a,b))}(1) = \text{Data-Loc}_{\text{SCM}}$.

(32) $\prod_{\text{AddressParts } \text{InsCode}(\text{MultBy}(a,b))}(2) = \text{Data-Loc}_{\text{SCM}}$.

(33) $\prod_{\text{AddressParts } \text{InsCode}(\text{Divide}(a,b))}(1) = \text{Data-Loc}_{\text{SCM}}$.

(34) $\prod_{\text{AddressParts } \text{InsCode}(\text{Divide}(a,b))}(2) = \text{Data-Loc}_{\text{SCM}}$.

(35) $\prod_{\text{AddressParts } \text{InsCode}(\text{goto } i_2)}(1) = \text{the instruction locations of } \text{SCM}$.

(36) $\prod_{\text{AddressParts } \text{InsCode}(\text{if } a=0 \text{ goto } i_2)}(1) = \text{the instruction locations of } \text{SCM}$.

(37) $\prod_{\text{AddressParts } \text{InsCode}(\text{if } a=0 \text{ goto } i_2)}(2) = \text{Data-Loc}_{\text{SCM}}$.

(38) $\prod_{\text{AddressParts } \text{InsCode}(\text{if } a>0 \text{ goto } i_2)}(1) = \text{the instruction locations of } \text{SCM}$.

(39) $\prod_{\text{AddressParts } \text{InsCode}(\text{if } a>0 \text{ goto } i_2)}(2) = \text{Data-Loc}_{\text{SCM}}$.

(40) $\text{NIC}(\text{halt}_{\text{SCM}}, i_1) = \{i_1\}$.

Let us note that $\text{JUMP}(\text{halt}_{\text{SCM}})$ is empty.

We now state the proposition

(41) $\text{NIC}(a:=b, i_1) = \{\text{Next}(i_1)\}$.

Let us consider a, b . Observe that $\text{JUMP}(a:=b)$ is empty.

We now state the proposition

(42) $\text{NIC}(\text{AddTo}(a,b), i_1) = \{\text{Next}(i_1)\}$.

Let us consider a, b . One can verify that $\text{JUMP}(\text{AddTo}(a,b))$ is empty.

We now state the proposition

$$(43) \quad \text{NIC}(\text{SubFrom}(a, b), i_1) = \{\text{Next}(i_1)\}.$$

Let us consider a, b . Observe that $\text{JUMP}(\text{SubFrom}(a, b))$ is empty.

We now state the proposition

$$(44) \quad \text{NIC}(\text{MultBy}(a, b), i_1) = \{\text{Next}(i_1)\}.$$

Let us consider a, b . Note that $\text{JUMP}(\text{MultBy}(a, b))$ is empty.

We now state the proposition

$$(45) \quad \text{NIC}(\text{Divide}(a, b), i_1) = \{\text{Next}(i_1)\}.$$

Let us consider a, b . One can verify that $\text{JUMP}(\text{Divide}(a, b))$ is empty.

One can prove the following propositions:

$$(46) \quad \text{NIC}(\text{goto } i_2, i_1) = \{i_2\}.$$

$$(47) \quad \text{JUMP}(\text{goto } i_2) = \{i_2\}.$$

Let us consider i_2 . One can verify that $\text{JUMP}(\text{goto } i_2)$ is non empty and trivial.

We now state two propositions:

$$(48) \quad \text{NIC}(\text{if } a = 0 \text{ goto } i_2, i_1) = \{i_2, \text{Next}(i_1)\}.$$

$$(49) \quad \text{JUMP}(\text{if } a = 0 \text{ goto } i_2) = \{i_2\}.$$

Let us consider a, i_2 . Observe that $\text{JUMP}(\text{if } a = 0 \text{ goto } i_2)$ is non empty and trivial.

One can prove the following propositions:

$$(50) \quad \text{NIC}(\text{if } a > 0 \text{ goto } i_2, i_1) = \{i_2, \text{Next}(i_1)\}.$$

$$(51) \quad \text{JUMP}(\text{if } a > 0 \text{ goto } i_2) = \{i_2\}.$$

Let us consider a, i_2 . Observe that $\text{JUMP}(\text{if } a > 0 \text{ goto } i_2)$ is non empty and trivial.

We now state two propositions:

$$(52) \quad \text{SUCC}(i_1) = \{i_1, \text{Next}(i_1)\}.$$

(53) Let f be a function from \mathbb{N} into the instruction locations of **SCM**. Suppose that for every natural number k holds $f(k) = \mathbf{i}_k$. Then

(i) f is bijective, and

(ii) for every natural number k holds $f(k+1) \in \text{SUCC}(f(k))$ and for every natural number j such that $f(j) \in \text{SUCC}(f(k))$ holds $k \leq j$.

Let us observe that **SCM** is standard.

We now state three propositions:

$$(54) \quad \text{il}_{\text{SCM}}(k) = \mathbf{i}_k.$$

$$(55) \quad \text{Next}(\text{il}_{\text{SCM}}(k)) = \text{il}_{\text{SCM}}(k+1).$$

$$(56) \quad \text{Next}(i_1) = \text{NextLoc } i_1.$$

Let us mention that $\text{InsCode}(\text{halt}_{\text{SCM}})$ is jump-only.

Let us note that halt_{SCM} is jump-only.

Let us consider i_2 . Note that $\text{InsCode}(\text{goto } i_2)$ is jump-only.

Let us consider i_2 . Observe that $\text{goto } i_2$ is jump-only, non sequential, and non instruction location free.

Let us consider a, i_2 . Note that $\text{InsCode}(\text{if } a = 0 \text{ goto } i_2)$ is jump-only and $\text{InsCode}(\text{if } a > 0 \text{ goto } i_2)$ is jump-only.

Let us consider a, i_2 . One can verify that $\text{if } a = 0 \text{ goto } i_2$ is jump-only, non sequential, and non instruction location free and $\text{if } a > 0 \text{ goto } i_2$ is jump-only, non sequential, and non instruction location free.

Let us consider a, b . One can verify the following observations:

- * $\text{InsCode}(a:=b)$ is non jump-only,
- * $\text{InsCode}(\text{AddTo}(a,b))$ is non jump-only,
- * $\text{InsCode}(\text{SubFrom}(a,b))$ is non jump-only,
- * $\text{InsCode}(\text{MultBy}(a,b))$ is non jump-only, and
- * $\text{InsCode}(\text{Divide}(a,b))$ is non jump-only.

Let us consider a, b . One can check the following observations:

- * $a:=b$ is non jump-only and sequential,
- * $\text{AddTo}(a,b)$ is non jump-only and sequential,
- * $\text{SubFrom}(a,b)$ is non jump-only and sequential,
- * $\text{MultBy}(a,b)$ is non jump-only and sequential, and
- * $\text{Divide}(a,b)$ is non jump-only and sequential.

Let us observe that **SCM** is homogeneous and has explicit jumps and no implicit jumps.

Let us note that **SCM** is regular.

We now state three propositions:

- (57) $\text{IncAddr}(\text{goto } i_2, k) = \text{goto } \text{ils}_{\text{SCM}}(\text{locnum}(i_2) + k)$.
- (58) $\text{IncAddr}(\text{if } a = 0 \text{ goto } i_2, k) = \text{if } a = 0 \text{ goto } \text{ils}_{\text{SCM}}(\text{locnum}(i_2) + k)$.
- (59) $\text{IncAddr}(\text{if } a > 0 \text{ goto } i_2, k) = \text{if } a > 0 \text{ goto } \text{ils}_{\text{SCM}}(\text{locnum}(i_2) + k)$.

Let us note that **SCM** is IC-good and Exec-preserving.

REFERENCES

- [1] Grzegorz Bancerek. The fundamental properties of natural numbers. *Journal of Formalized Mathematics*, 1, 1989. http://mizar.org/JFM/Vol1/nat_1.html.
- [2] Grzegorz Bancerek. The ordinal numbers. *Journal of Formalized Mathematics*, 1, 1989. <http://mizar.org/JFM/Vol1/ordinal1.html>.
- [3] Grzegorz Bancerek. Sequences of ordinal numbers. *Journal of Formalized Mathematics*, 1, 1989. <http://mizar.org/JFM/Vol1/ordinal2.html>.
- [4] Grzegorz Bancerek. König's theorem. *Journal of Formalized Mathematics*, 2, 1990. http://mizar.org/JFM/Vol2/card_3.html.
- [5] Grzegorz Bancerek and Krzysztof Hryniewiecki. Segments of natural numbers and finite sequences. *Journal of Formalized Mathematics*, 1, 1989. http://mizar.org/JFM/Vol1/finseq_1.html.
- [6] Grzegorz Bancerek and Andrzej Trybulec. Miscellaneous facts about functions. *Journal of Formalized Mathematics*, 8, 1996. http://mizar.org/JFM/Vol8/funct_7.html.
- [7] Czesław Byliński. Functions and their basic properties. *Journal of Formalized Mathematics*, 1, 1989. http://mizar.org/JFM/Vol1/funct_1.html.
- [8] Czesław Byliński. Functions from a set to a set. *Journal of Formalized Mathematics*, 1, 1989. http://mizar.org/JFM/Vol1/funct_2.html.
- [9] Czesław Byliński. A classical first order language. *Journal of Formalized Mathematics*, 2, 1990. http://mizar.org/JFM/Vol2/cqc_lang.html.
- [10] Czesław Byliński. The modification of a function by a function and the iteration of the composition of a function. *Journal of Formalized Mathematics*, 2, 1990. http://mizar.org/JFM/Vol2/funct_4.html.
- [11] Artur Kornilowicz. On the composition of macro instructions of standard computers. *Journal of Formalized Mathematics*, 12, 2000. http://mizar.org/JFM/Vol12/amistd_2.html.
- [12] Yatsuka Nakamura and Andrzej Trybulec. A mathematical model of CPU. *Journal of Formalized Mathematics*, 4, 1992. http://mizar.org/JFM/Vol4/ami_1.html.
- [13] Yatsuka Nakamura and Andrzej Trybulec. On a mathematical model of programs. *Journal of Formalized Mathematics*, 4, 1992. http://mizar.org/JFM/Vol4/ami_2.html.

- [14] Beata Padlewska. Families of sets. *Journal of Formalized Mathematics*, 1, 1989. http://mizar.org/JFM/Vol1/setfam_1.html.
- [15] Dariusz Surowik. Cyclic groups and some of their properties — part I. *Journal of Formalized Mathematics*, 3, 1991. http://mizar.org/JFM/Vol3/gr_cy_1.html.
- [16] Yasushi Tanaka. On the decomposition of the states of SCM. *Journal of Formalized Mathematics*, 5, 1993. http://mizar.org/JFM/Vol5/ami_5.html.
- [17] Andrzej Trybulec. Tarski Grothendieck set theory. *Journal of Formalized Mathematics*, Axiomatics, 1989. <http://mizar.org/JFM/Axiomatics/tarski.html>.
- [18] Andrzej Trybulec. Tuples, projections and Cartesian products. *Journal of Formalized Mathematics*, 1, 1989. http://mizar.org/JFM/Vol1/mcart_1.html.
- [19] Andrzej Trybulec. Subsets of real numbers. *Journal of Formalized Mathematics*, Addenda, 2003. <http://mizar.org/JFM/Addenda/numbers.html>.
- [20] Andrzej Trybulec and Yatsuka Nakamura. Some remarks on the simple concrete model of computer. *Journal of Formalized Mathematics*, 5, 1993. http://mizar.org/JFM/Vol5/ami_3.html.
- [21] Andrzej Trybulec, Piotr Rudnicki, and Artur Korniłowicz. Standard ordering of instruction locations. *Journal of Formalized Mathematics*, 12, 2000. http://mizar.org/JFM/Vol12/amistd_1.html.
- [22] Michał J. Trybulec. Integers. *Journal of Formalized Mathematics*, 2, 1990. http://mizar.org/JFM/Vol2/int_1.html.
- [23] Zinaida Trybulec. Properties of subsets. *Journal of Formalized Mathematics*, 1, 1989. http://mizar.org/JFM/Vol1/subset_1.html.
- [24] Edmund Woronowicz. Relations and their basic properties. *Journal of Formalized Mathematics*, 1, 1989. http://mizar.org/JFM/Vol1/relat_1.html.

Received May 8, 2001

Published January 2, 2004
