

# On a Mathematical Model of Programs

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**Summary.** We continue the work on mathematical modeling of hardware and software started in [11]. The main objective of this paper is the definition of a program. We start with the concept of partial product, i.e. the set of all partial functions  $f$  from  $I$  to  $\bigcup_{i \in I} A_i$ , fulfilling the condition  $f.i \in A_i$  for  $i \in \text{dom}f$ . The computation and the result of a computation are defined in usual way. A finite partial state is called autonomic if the result of a computation starting with it does not depend on the remaining memory and an AMI is called programmable if it has a non empty autonomic partial finite state. We prove the consistency of the following set of properties of an AMI: data-oriented, halting, steady-programmed, realistic and programmable. For this purpose we define a trivial AMI. It has only the instruction counter and one instruction location. The only instruction of it is the halt instruction. A pre-program is a finite partial state that halts. We conclude with the definition of a program of a partial function  $F$  mapping the set of the finite partial states into itself. It is a finite partial state  $s$  such that for every finite partial state  $s' \in \text{dom}F$  the result of any computation starting with  $s + s'$  includes  $F.s'$ .

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The articles [15], [14], [7], [20], [2], [17], [3], [21], [5], [6], [12], [13], [18], [1], [8], [16], [9], [10], [4], and [19] provide the notation and terminology for this paper.

In this paper  $x$  denotes a set and  $i, k$  denote natural numbers.

The element  $\text{Halt}_{\text{SCM}}$  of  $\mathbb{Z}_9$  is defined as follows:

(Def. 1)  $\text{Halt}_{\text{SCM}} = 0$ .

The subset  $\text{Data-Loc}_{\text{SCM}}$  of  $\mathbb{N}$  is defined by:

(Def. 2)  $\text{Data-Loc}_{\text{SCM}} = \{2 \cdot k + 1\}$ .

The subset  $\text{Instr-Loc}_{\text{SCM}}$  of  $\mathbb{N}$  is defined by:

(Def. 3)  $\text{Instr-Loc}_{\text{SCM}} = \{2 \cdot k : k > 0\}$ .

Let us observe that  $\text{Data-Loc}_{\text{SCM}}$  is non empty and  $\text{Instr-Loc}_{\text{SCM}}$  is non empty.

We use the following convention:  $I, J, K$  denote elements of  $\mathbb{Z}_9$ ,  $a, a_1, a_2$  denote elements of  $\text{Instr-Loc}_{\text{SCM}}$ , and  $b, b_1, b_2, c, c_1$  denote elements of  $\text{Data-Loc}_{\text{SCM}}$ .

The subset  $\text{Instr}_{\text{SCM}}$  of  $[\mathbb{Z}_9, (\bigcup\{\mathbb{Z}\} \cup \mathbb{N})^*]$  is defined by:

(Def. 4)  $\text{Instr}_{\text{SCM}} = \{\langle \text{Halt}_{\text{SCM}}, \emptyset \rangle\} \cup \{\langle J, \langle a \rangle \rangle : J = 6\} \cup \{\langle K, \langle a_1, b_1 \rangle \rangle : K \in \{7, 8\}\} \cup \{\langle I, \langle b, c \rangle \rangle : I \in \{1, 2, 3, 4, 5\}\}$ .

The following proposition is true

(2)<sup>1</sup>  $\langle 0, \emptyset \rangle \in \text{Instr}_{\text{SCM}}$ .

One can check that  $\text{Instr}_{\text{SCM}}$  is non empty.

One can prove the following three propositions:

(3)  $\langle 6, \langle a_2 \rangle \rangle \in \text{Instr}_{\text{SCM}}$ .

(4) If  $x \in \{7, 8\}$ , then  $\langle x, \langle a_2, b_2 \rangle \rangle \in \text{Instr}_{\text{SCM}}$ .

(5) If  $x \in \{1, 2, 3, 4, 5\}$ , then  $\langle x, \langle b_1, c_1 \rangle \rangle \in \text{Instr}_{\text{SCM}}$ .

The function  $\text{OK}_{\text{SCM}}$  from  $\mathbb{N}$  into  $\{\mathbb{Z}\} \cup \{\text{Instr}_{\text{SCM}}, \text{Instr-Loc}_{\text{SCM}}\}$  is defined by:

(Def. 5)  $\text{OK}_{\text{SCM}}(0) = \text{Instr-Loc}_{\text{SCM}}$  and for every natural number  $k$  holds  $\text{OK}_{\text{SCM}}(2 \cdot k + 1) = \mathbb{Z}$  and  $\text{OK}_{\text{SCM}}(2 \cdot k + 2) = \text{Instr}_{\text{SCM}}$ .

One can prove the following propositions:

(6)  $\text{Instr-Loc}_{\text{SCM}} \neq \mathbb{Z}$  and  $\text{Instr}_{\text{SCM}} \neq \mathbb{Z}$  and  $\text{Instr-Loc}_{\text{SCM}} \neq \text{Instr}_{\text{SCM}}$ .

(7)  $\text{OK}_{\text{SCM}}(i) = \text{Instr-Loc}_{\text{SCM}}$  iff  $i = 0$ .

(8)  $\text{OK}_{\text{SCM}}(i) = \mathbb{Z}$  iff there exists  $k$  such that  $i = 2 \cdot k + 1$ .

(9)  $\text{OK}_{\text{SCM}}(i) = \text{Instr}_{\text{SCM}}$  iff there exists  $k$  such that  $i = 2 \cdot k + 2$ .

A **SCM**-state is an element of  $\prod(\text{OK}_{\text{SCM}})$ .

Next we state several propositions:

(10) For every element  $a$  of  $\text{Data-Loc}_{\text{SCM}}$  holds  $\text{OK}_{\text{SCM}}(a) = \mathbb{Z}$ .

(11) For every element  $a$  of  $\text{Instr-Loc}_{\text{SCM}}$  holds  $\text{OK}_{\text{SCM}}(a) = \text{Instr}_{\text{SCM}}$ .

(12) For every element  $a$  of  $\text{Instr-Loc}_{\text{SCM}}$  and for every element  $t$  of  $\text{Data-Loc}_{\text{SCM}}$  holds  $a \neq t$ .

(13)  $\pi_0 \prod(\text{OK}_{\text{SCM}}) = \text{Instr-Loc}_{\text{SCM}}$ .

(14) For every element  $a$  of  $\text{Data-Loc}_{\text{SCM}}$  holds  $\pi_a \prod(\text{OK}_{\text{SCM}}) = \mathbb{Z}$ .

(15) For every element  $a$  of  $\text{Instr-Loc}_{\text{SCM}}$  holds  $\pi_a \prod(\text{OK}_{\text{SCM}}) = \text{Instr}_{\text{SCM}}$ .

Let  $s$  be a **SCM**-state. The functor  $\mathbf{IC}_s$  yielding an element of  $\text{Instr-Loc}_{\text{SCM}}$  is defined by:

(Def. 6)  $\mathbf{IC}_s = s(0)$ .

Let  $s$  be a **SCM**-state and let  $u$  be an element of  $\text{Instr-Loc}_{\text{SCM}}$ . The functor  $\text{Chg}_{\text{SCM}}(s, u)$  yields a **SCM**-state and is defined as follows:

(Def. 7)  $\text{Chg}_{\text{SCM}}(s, u) = s + \cdot (0 \dashrightarrow u)$ .

We now state three propositions:

(16) For every **SCM**-state  $s$  and for every element  $u$  of  $\text{Instr-Loc}_{\text{SCM}}$  holds  $(\text{Chg}_{\text{SCM}}(s, u))(0) = u$ .

(17) For every **SCM**-state  $s$  and for every element  $u$  of  $\text{Instr-Loc}_{\text{SCM}}$  and for every element  $m_1$  of  $\text{Data-Loc}_{\text{SCM}}$  holds  $(\text{Chg}_{\text{SCM}}(s, u))(m_1) = s(m_1)$ .

(18) For every **SCM**-state  $s$  and for all elements  $u, v$  of  $\text{Instr-Loc}_{\text{SCM}}$  holds  $(\text{Chg}_{\text{SCM}}(s, u))(v) = s(v)$ .

Let  $s$  be a **SCM**-state, let  $t$  be an element of  $\text{Data-Loc}_{\text{SCM}}$ , and let  $u$  be an integer. The functor  $\text{Chg}_{\text{SCM}}(s, t, u)$  yields a **SCM**-state and is defined by:

<sup>1</sup> The proposition (1) has been removed.

(Def. 8)  $\text{Chg}_{\text{SCM}}(s, t, u) = s + \cdot (t \mapsto u)$ .

Next we state four propositions:

(19) For every **SCM**-state  $s$  and for every element  $t$  of  $\text{Data-Loc}_{\text{SCM}}$  and for every integer  $u$  holds  $(\text{Chg}_{\text{SCM}}(s, t, u))(0) = s(0)$ .

(20) For every **SCM**-state  $s$  and for every element  $t$  of  $\text{Data-Loc}_{\text{SCM}}$  and for every integer  $u$  holds  $(\text{Chg}_{\text{SCM}}(s, t, u))(t) = u$ .

(21) Let  $s$  be a **SCM**-state,  $t$  be an element of  $\text{Data-Loc}_{\text{SCM}}$ ,  $u$  be an integer, and  $m_1$  be an element of  $\text{Data-Loc}_{\text{SCM}}$ . If  $m_1 \neq t$ , then  $(\text{Chg}_{\text{SCM}}(s, t, u))(m_1) = s(m_1)$ .

(22) Let  $s$  be a **SCM**-state,  $t$  be an element of  $\text{Data-Loc}_{\text{SCM}}$ ,  $u$  be an integer, and  $v$  be an element of  $\text{Instr-Loc}_{\text{SCM}}$ . Then  $(\text{Chg}_{\text{SCM}}(s, t, u))(v) = s(v)$ .

Let  $x$  be an element of  $\text{Instr}_{\text{SCM}}$ . Let us assume that there exist elements  $m_1, m_2$  of  $\text{Data-Loc}_{\text{SCM}}$  and  $I$  such that  $x = \langle I, \langle m_1, m_2 \rangle \rangle$ . The functor  $x\text{address}_1$  yielding an element of  $\text{Data-Loc}_{\text{SCM}}$  is defined as follows:

(Def. 9) There exists a finite sequence  $f$  of elements of  $\text{Data-Loc}_{\text{SCM}}$  such that  $f = x_2$  and  $x\text{address}_1 = f_1$ .

The functor  $x\text{address}_2$  yields an element of  $\text{Data-Loc}_{\text{SCM}}$  and is defined by:

(Def. 10) There exists a finite sequence  $f$  of elements of  $\text{Data-Loc}_{\text{SCM}}$  such that  $f = x_2$  and  $x\text{address}_2 = f_2$ .

One can prove the following proposition

(23) For every element  $x$  of  $\text{Instr}_{\text{SCM}}$  and for all elements  $m_1, m_2$  of  $\text{Data-Loc}_{\text{SCM}}$  and for every  $I$  such that  $x = \langle I, \langle m_1, m_2 \rangle \rangle$  holds  $x\text{address}_1 = m_1$  and  $x\text{address}_2 = m_2$ .

Let  $x$  be an element of  $\text{Instr}_{\text{SCM}}$ . Let us assume that there exist an element  $m_1$  of  $\text{Instr-Loc}_{\text{SCM}}$  and  $I$  such that  $x = \langle I, \langle m_1 \rangle \rangle$ . The functor  $x\text{address}_j$  yields an element of  $\text{Instr-Loc}_{\text{SCM}}$  and is defined as follows:

(Def. 11) There exists a finite sequence  $f$  of elements of  $\text{Instr-Loc}_{\text{SCM}}$  such that  $f = x_2$  and  $x\text{address}_j = f_1$ .

One can prove the following proposition

(24) For every element  $x$  of  $\text{Instr}_{\text{SCM}}$  and for every element  $m_1$  of  $\text{Instr-Loc}_{\text{SCM}}$  and for every  $I$  such that  $x = \langle I, \langle m_1 \rangle \rangle$  holds  $x\text{address}_j = m_1$ .

Let  $x$  be an element of  $\text{Instr}_{\text{SCM}}$ . Let us assume that there exist an element  $m_1$  of  $\text{Instr-Loc}_{\text{SCM}}$ , an element  $m_2$  of  $\text{Data-Loc}_{\text{SCM}}$ , and  $I$  such that  $x = \langle I, \langle m_1, m_2 \rangle \rangle$ . The functor  $x\text{address}_j$  yields an element of  $\text{Instr-Loc}_{\text{SCM}}$  and is defined by:

(Def. 12) There exists an element  $m_1$  of  $\text{Instr-Loc}_{\text{SCM}}$  and there exists an element  $m_2$  of  $\text{Data-Loc}_{\text{SCM}}$  such that  $\langle m_1, m_2 \rangle = x_2$  and  $x\text{address}_j = \langle m_1, m_2 \rangle_1$ .

The functor  $x\text{address}_c$  yielding an element of  $\text{Data-Loc}_{\text{SCM}}$  is defined by:

(Def. 13) There exists an element  $m_1$  of  $\text{Instr-Loc}_{\text{SCM}}$  and there exists an element  $m_2$  of  $\text{Data-Loc}_{\text{SCM}}$  such that  $\langle m_1, m_2 \rangle = x_2$  and  $x\text{address}_c = \langle m_1, m_2 \rangle_2$ .

Next we state the proposition

(25) Let  $x$  be an element of  $\text{Instr}_{\text{SCM}}$ ,  $m_1$  be an element of  $\text{Instr-Loc}_{\text{SCM}}$ ,  $m_2$  be an element of  $\text{Data-Loc}_{\text{SCM}}$ , and given  $I$ . If  $x = \langle I, \langle m_1, m_2 \rangle \rangle$ , then  $x\text{address}_j = m_1$  and  $x\text{address}_c = m_2$ .

Let  $s$  be a **SCM**-state and let  $a$  be an element of  $\text{Data-Loc}_{\text{SCM}}$ . Note that  $s(a)$  is integer.

Let  $D$  be a non empty set, let  $x, y$  be real numbers, and let  $a, b$  be elements of  $D$ . The functor  $(x > y \rightarrow a, b)$  yields an element of  $D$  and is defined as follows:

$$\text{(Def. 14)} \quad (x > y \rightarrow a, b) = \begin{cases} a, & \text{if } x > y, \\ b, & \text{otherwise.} \end{cases}$$

Let  $d$  be an element of  $\text{Instr-Loc}_{\text{SCM}}$ . The functor  $\text{Next}(d)$  yields an element of  $\text{Instr-Loc}_{\text{SCM}}$  and is defined by:

$$\text{(Def. 15)} \quad \text{Next}(d) = d + 2.$$

Let  $x$  be an element of  $\text{Instr}_{\text{SCM}}$  and let  $s$  be a **SCM**-state. The functor  $\text{Exec-Res}_{\text{SCM}}(x, s)$  yields a **SCM**-state and is defined as follows:

$$\text{(Def. 16)} \quad \text{Exec-Res}_{\text{SCM}}(x, s) = \begin{cases} \text{Chg}_{\text{SCM}}(\text{Chg}_{\text{SCM}}(s, x \text{ address}_1, s(x \text{ address}_2)), \text{Next}(\mathbf{IC}_s)), & \text{if there exist elements } m_1, m_2 \text{ of } \text{Instr-Loc}_{\text{SCM}} \text{ such that } x = \langle 6, \langle m_1, m_2 \rangle \rangle, \\ \text{Chg}_{\text{SCM}}(\text{Chg}_{\text{SCM}}(s, x \text{ address}_1, s(x \text{ address}_1) + s(x \text{ address}_2)), \text{Next}(\mathbf{IC}_s)), & \text{if there exist elements } m_1, m_2 \text{ of } \text{Instr-Loc}_{\text{SCM}} \text{ such that } x = \langle 7, \langle m_1, m_2 \rangle \rangle, \\ \text{Chg}_{\text{SCM}}(\text{Chg}_{\text{SCM}}(s, x \text{ address}_1, s(x \text{ address}_1) - s(x \text{ address}_2)), \text{Next}(\mathbf{IC}_s)), & \text{if there exist elements } m_1, m_2 \text{ of } \text{Instr-Loc}_{\text{SCM}} \text{ such that } x = \langle 8, \langle m_1, m_2 \rangle \rangle, \\ \text{Chg}_{\text{SCM}}(\text{Chg}_{\text{SCM}}(s, x \text{ address}_1, s(x \text{ address}_1) \cdot s(x \text{ address}_2)), \text{Next}(\mathbf{IC}_s)), & \text{if there exist elements } m_1, m_2 \text{ of } \text{Instr-Loc}_{\text{SCM}} \text{ such that } x = \langle 9, \langle m_1, m_2 \rangle \rangle, \\ \text{Chg}_{\text{SCM}}(\text{Chg}_{\text{SCM}}(\text{Chg}_{\text{SCM}}(s, x \text{ address}_1, s(x \text{ address}_1) \div s(x \text{ address}_2)), x \text{ address}_2, s(x \text{ address}_2)), & \text{if there exist elements } m_1, m_2 \text{ of } \text{Instr-Loc}_{\text{SCM}} \text{ such that } x = \langle 10, \langle m_1, m_2 \rangle \rangle, \\ \text{Chg}_{\text{SCM}}(s, x \text{ address}_j), & \text{if there exists an element } m_1 \text{ of } \text{Instr-Loc}_{\text{SCM}} \text{ such that } x = \langle 6, \langle m_1 \rangle \rangle, \\ \text{Chg}_{\text{SCM}}(s, (s(x \text{ address}_c) = 0 \rightarrow x \text{ address}_j, \text{Next}(\mathbf{IC}_s))), & \text{if there exists an element } m_1 \text{ of } \text{Instr-Loc}_{\text{SCM}} \text{ such that } x = \langle 11, m_1 \rangle, \\ \text{Chg}_{\text{SCM}}(s, (s(x \text{ address}_c) > 0 \rightarrow x \text{ address}_j, \text{Next}(\mathbf{IC}_s))), & \text{if there exists an element } m_1 \text{ of } \text{Instr-Loc}_{\text{SCM}} \text{ such that } x = \langle 12, m_1 \rangle, \\ s, & \text{otherwise.} \end{cases}$$

The function  $\text{Exec}_{\text{SCM}}$  from  $\text{Instr}_{\text{SCM}}$  into  $(\prod(\text{OK}_{\text{SCM}}))^{\prod(\text{OK}_{\text{SCM}})}$  is defined by:

$$\text{(Def. 17)} \quad \text{For every element } x \text{ of } \text{Instr}_{\text{SCM}} \text{ and for every } \text{SCM}\text{-state } y \text{ holds } \text{Exec}_{\text{SCM}}(x)(y) = \text{Exec-Res}_{\text{SCM}}(x, y).$$

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